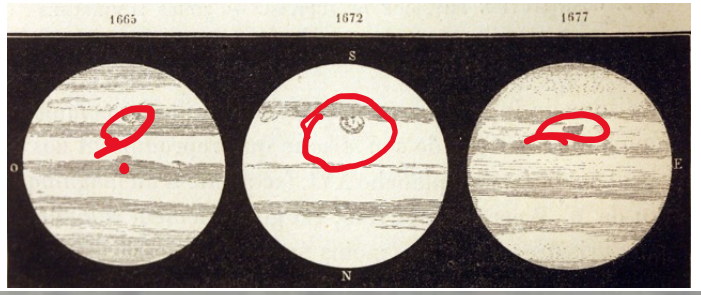
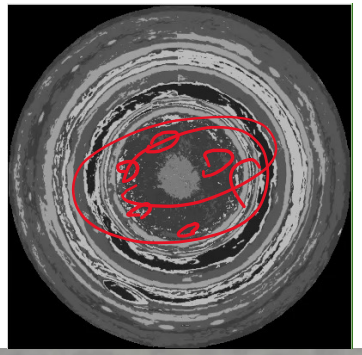
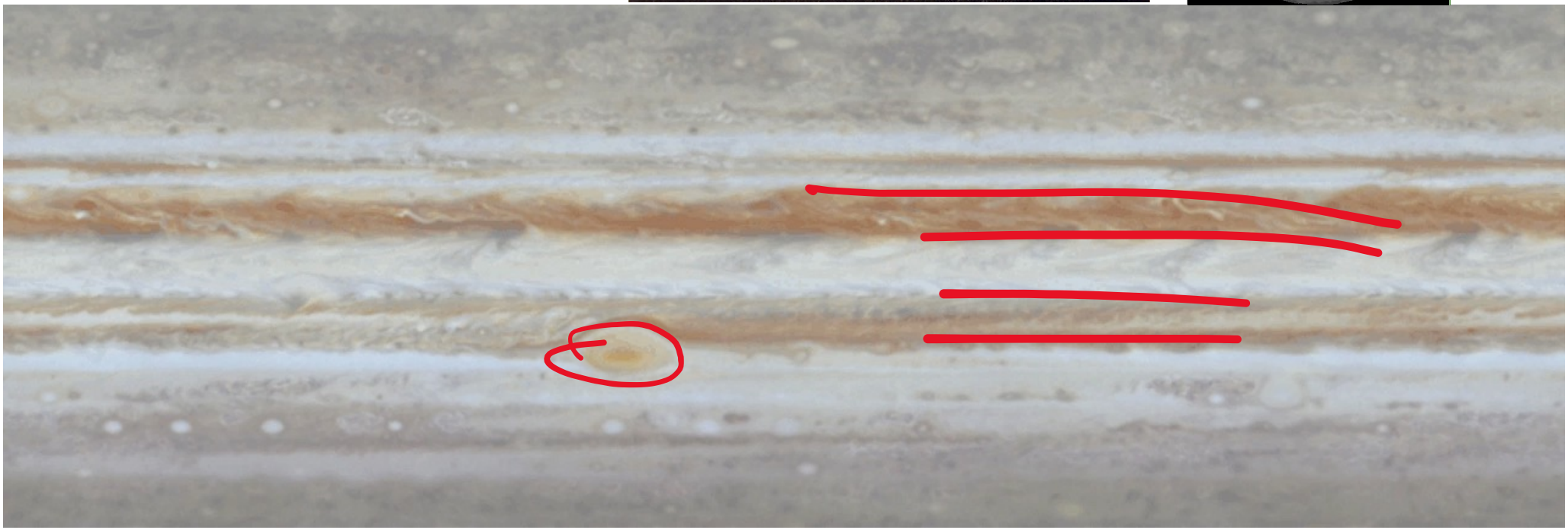
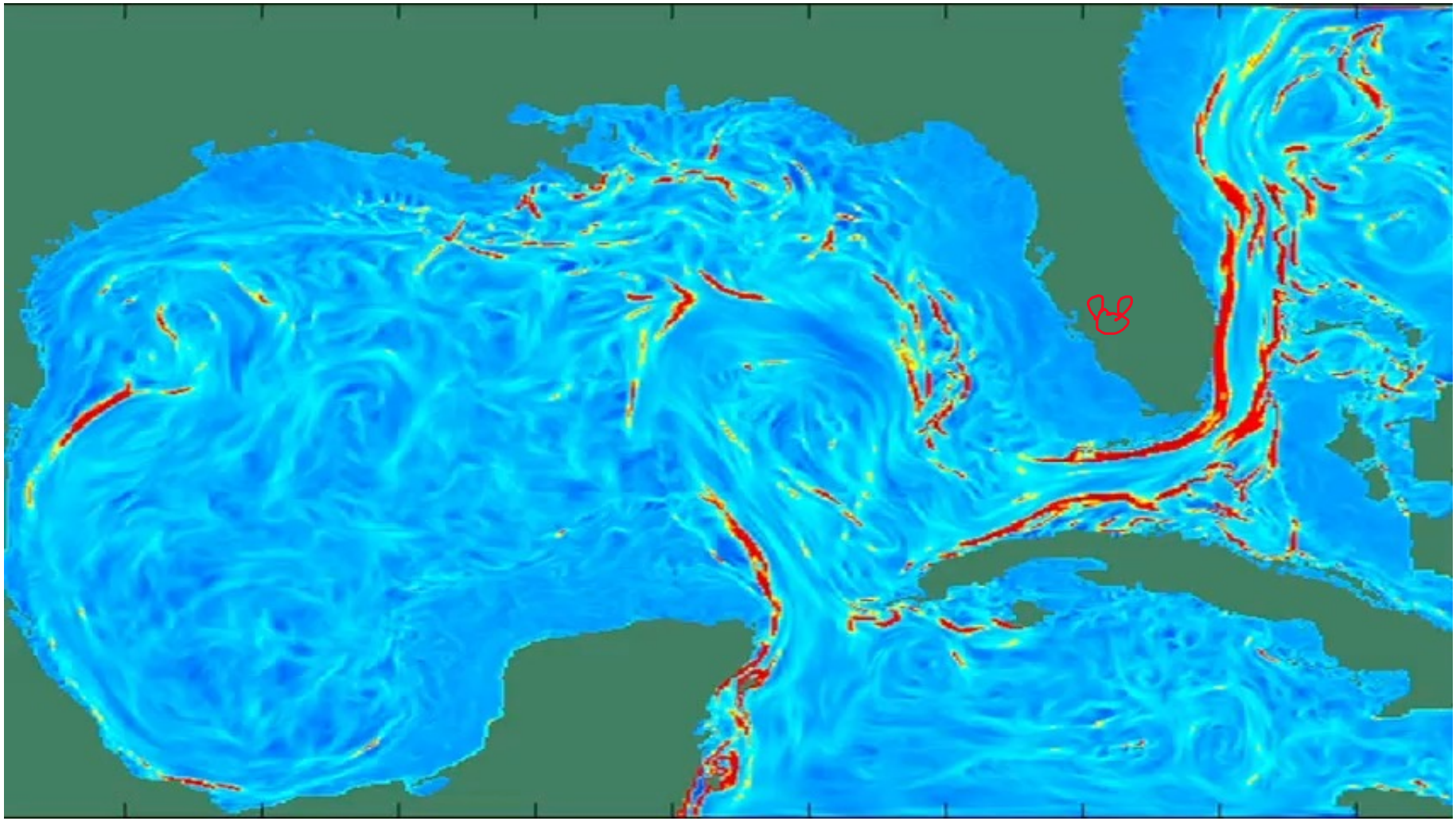
EE520 Data Driven Analysis of Complex Systems

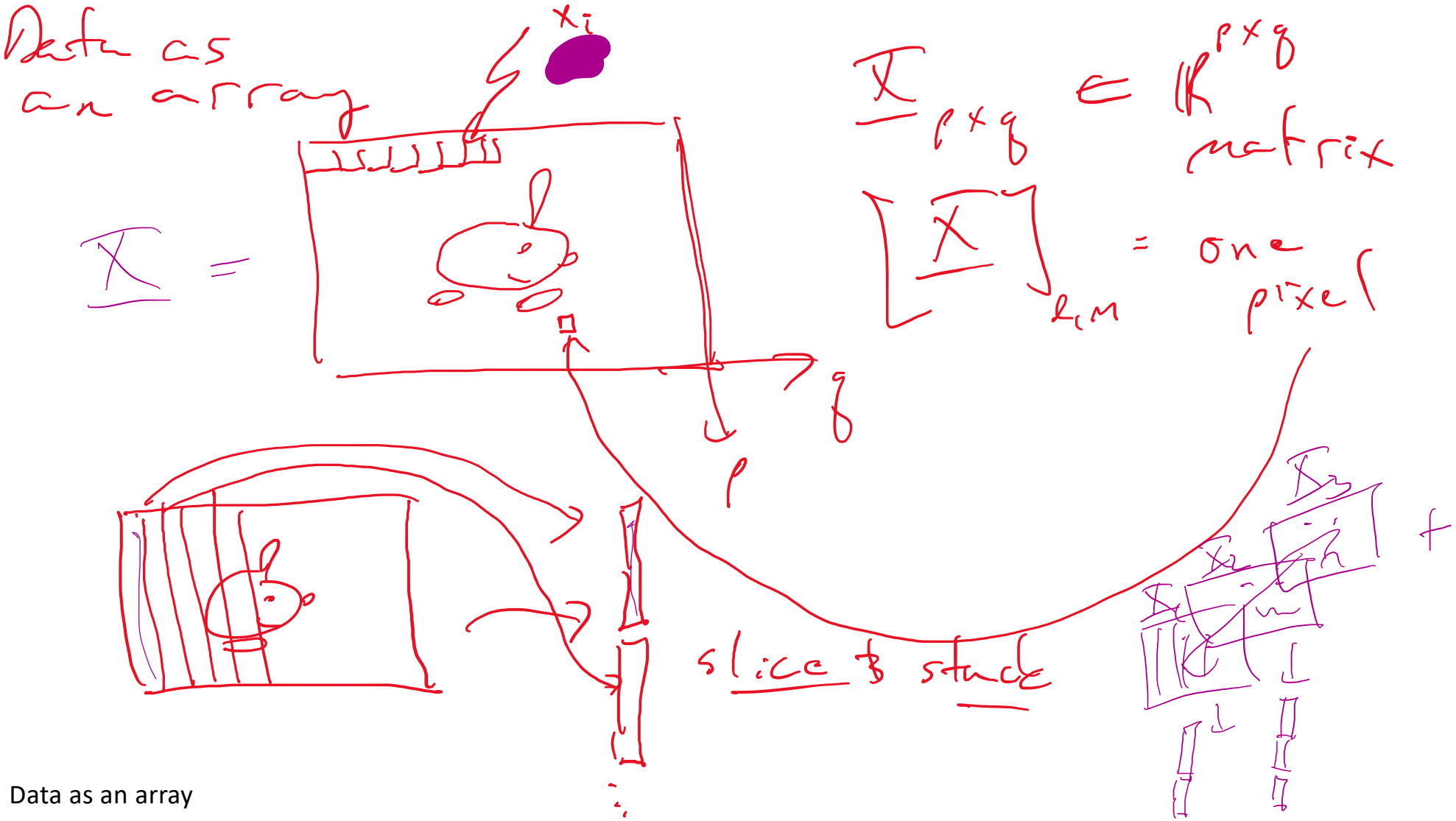
Erik Bollt











Data as an array

On Matrix Multiplication

$$L(\mathbf{z}) : \mathbb{R}^n \to \mathbb{R}^m, \quad \mathbf{z} \mapsto \mathbf{z}' = A\mathbf{z},$$

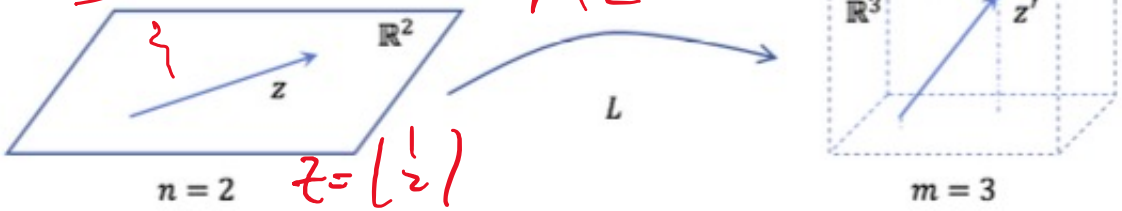
in terms of the usual matrix times vector multiplication,

$$A = \begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n} \\ a_{2,1} & a_{2,2} & \cdots & a_{2,n} \\ \vdots & & \ddots & \vdots \\ a_{m,1} & a_{m,2} & \cdots & a_{m,n} \end{pmatrix}, \text{ and each } a_{i,j} \in \mathbb{C}$$

$$[\mathbf{z}]'_i = \sum_{j=1}^{n} A_{i,j}[\mathbf{z}]_j, \text{ for each } i = 1, \cdots, m,$$

- a vector has length & direction.

Geometric



$Az = z'$

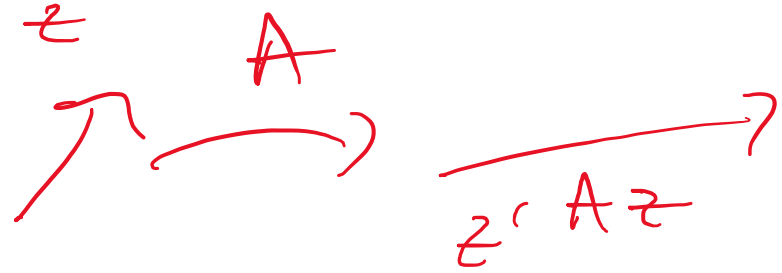
$z = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$

But as linear algebra

$A_{2\times 2} = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}$; matrice × vectors

$A\begin{pmatrix} 3 \\ 4 \end{pmatrix} = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}\begin{pmatrix} 3 \\ 4 \end{pmatrix} = \begin{pmatrix} 1\cdot 3 + 2\cdot 4 \\ 3\cdot 3 + 4\cdot 4 \end{pmatrix} = \begin{pmatrix} 11 \\ 25 \end{pmatrix}$

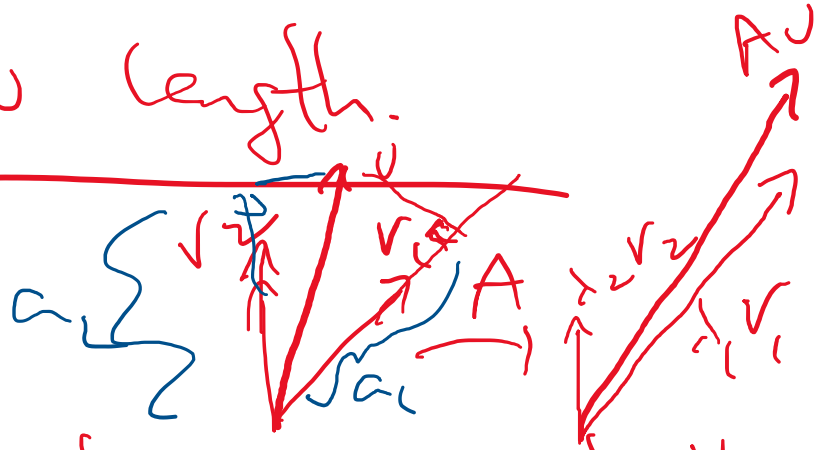
$z' = Az$



new direction, new length.

Eig for square.

- $Av = \lambda v$
  2×2



Characterize matrices by knowing just these

- e.g. special directions

- $Au = A(a_1 v_1 + a_2 v_2) = a_1 A v_1 + a_2 A v_2$
  $= a_1 \lambda_1 v_1 + a_2 \lambda_2 v_2$

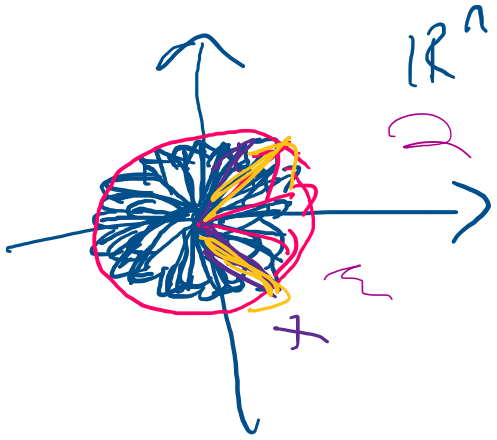
$\det(A - \lambda I) = 0$

$(A - \lambda I)v = 0$

Matrix time circle =

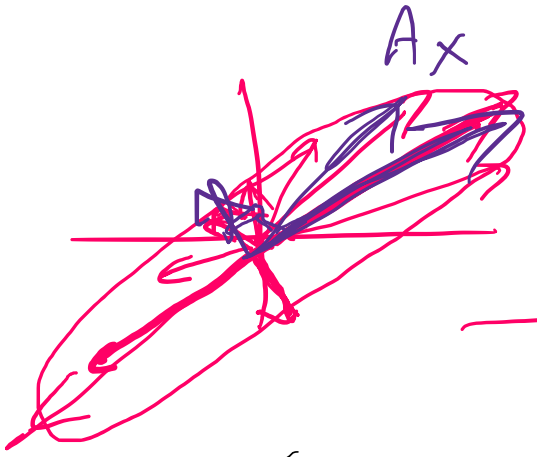
all vectors of length 1.

? Matrix x circle ?! But matrix times vector. =



A

Ax



$S = \{x \mid \|x\|_2 = 1, x \in E \cong \mathbb{R}^2\}$; $A \cdot S = \{y : y = Ax, x \in S\}$

Theorem 2.1.1 — Singular Value Decomposition. Let A be an $m \times n$ matrix whose entries come from the field $\mathcal{K}$, which is either the field of real numbers or the field of complex numbers. Then the singular value decomposition of A exists, and it takes the form of a product of matrices:

$$A_{m\times n} = U_{m\times m}\Sigma_{m\times n}V^*_{n\times n}, \tag{2.5}$$

where

- U is an $m \times m$ unitary matrix.
- Σ is a diagonal $m \times n$ matrix with non-negative real numbers on the diagonal.
- V is an $n \times n$ unitary matrix, and V^* is the conjugate transpose of V.

$$\Sigma := diag(\sigma_1, \sigma_2, \cdots, \sigma_p), p = min(m, n),$$
$$\sigma_1 \geq \sigma_2 \geq \cdots \geq \sigma_p \geq 0.$$

$$A = U \Sigma V^T$$

The singular values are the nonegative values: $\sigma_i \geq 0, i = 1, \cdots, n,$

The left singular vectors: u_i are the columns of $U = [u_1, u_2, ..., u_m]$.

The right singular vectors: v_i are the columns of $V = [v_1, v_2, ..., v_n]$

Definition 2.1.1 — Singular values and singular vectors. The singular values of A are the scalar values, σ_i, and the columns of U and V have columns that are the corresponding i^{th} left and right singular vectors, u_i and v_i:

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Since V is orthogonal, then right multiplying Eq. (2.5) by V,

$$AV = U\Sigma V^*V = U\Sigma, \tag{2.8}$$

Fri 08/21/20

$$Ax = b$$

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 6 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

if A is square

$$A\,V = U\,\Sigma$$

$$[A]\,[v_1 v_2 \cdots v_n] = [u_1 u_2 \cdots u_n]\, diag(\sigma_1, \sigma_2, \cdots, \sigma_n).$$

■ **Example 2.1** Let $A = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 6 \end{pmatrix}_{2\times 3}$. By SVD of the matrix A we have:

$$\begin{aligned} A &= U\Sigma V^T \\ &= \begin{pmatrix} \frac{1}{\sqrt{5}} & \frac{2}{\sqrt{5}} \\ \frac{-2}{\sqrt{5}} & \frac{1}{\sqrt{5}} \end{pmatrix} \begin{pmatrix} \sqrt{70} & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{14}} & \sqrt{\frac{2}{7}} & \frac{3}{\sqrt{14}} \\ \frac{-3}{\sqrt{10}} & 0 & \frac{1}{\sqrt{10}} \\ \frac{-1}{\sqrt{35}} & \sqrt{\frac{5}{7}} & \frac{-3}{\sqrt{35}} \end{pmatrix} \quad (2.28) \end{aligned}$$

We see that the second singular value, $\sigma_2 = 2$, meaning that number of non-zero singular values $r < \min\{m, n\}$. Such matrix is called rank deficient matrix. If we take the economy version (with $r = 1$) of the SVD we will have:

$$\begin{aligned} u_1\sigma_1 v_1^T &= \begin{pmatrix} \frac{1}{\sqrt{5}} \\ \frac{-2}{\sqrt{5}} \end{pmatrix} \begin{pmatrix} \sqrt{70} \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{14}} & \sqrt{\frac{2}{7}} & \frac{3}{\sqrt{14}} \end{pmatrix} \\ &\approx \begin{pmatrix} 1 & 2 & 3 \\ 2 & 4 & 6 \end{pmatrix} \quad (2.29) \end{aligned}$$

$$A^T A = V \Sigma^T \Sigma V^T \cdot V$$

$$A^T A V = \Sigma^T \Sigma V$$

$$V = \begin{pmatrix} v_1 & v_2 & \cdots & v_n \end{pmatrix}$$

$$[A]\,[v_1 v_2 \cdots v_n] = [u_1 u_2 \cdots u_n]\, diag(\sigma_1, \sigma_2, \cdots, \sigma_n).$$

Full ✱

The Economy SVD, and Reduced Rank SVD

The general SVD, Eq. (2.5) may be written in terms of submatrices.

Definition 2.1.3 — The Economy SVD. For any matrix $A \in \mathbb{R}^{m\times n}$, the general SVD Eq. (2.5) can be written in terms of smaller matrices,

$$A_{m\times n} = \hat{U}_{m\times n}\hat{\Sigma}_{n\times n}V^*_{n\times n}, \tag{2.21}$$

and $U = [\hat{U}_{m\times n}|\hat{U}_{(n-m)\times n}]$, written in terms of an orthogonal "buffer" matrix

$\sigma_1 \geq \sigma_2 \geq \ldots \geq \sigma_r > \sigma_{r+1} = 0$

$= 0 \ldots$

Definition 2.1.4 — Rank Deficient SVD. For a matrix $A \in \mathbb{R}^{m\times n}$ such that the SVD results in singular values

$$\sigma_r > \sigma_{r+1} = 0, \text{ for some } r < n. \tag{2.22}$$

then the SVD can be written in terms of an economy form as smaller matrices,

$$A_{m\times n} = \hat{U}_{m\times r}\hat{\Sigma}_{n\times n}V^*_{n\times r}, \tag{2.23}$$

and related to the general SVD Eq. (2.5) by $U = [\hat{U}_{m\times r}|\hat{U}_{(n-r)\times n}]$, but $r < n$.

Rank ill conditioned

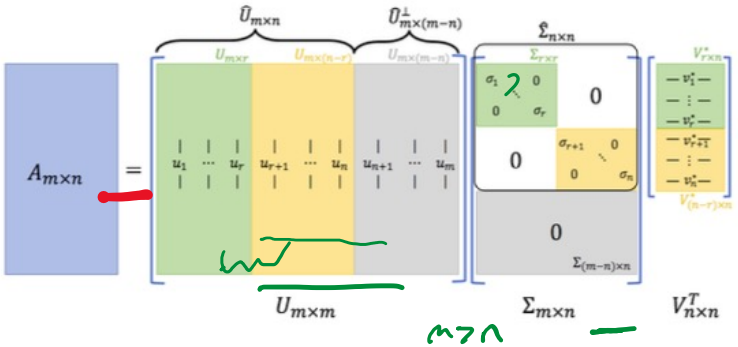


Figure 2.3: $m > n$ tall skinny

Recall that,

$$A_{m\times n} = \hat{U}_{m\times n}\hat{\Sigma}_{n\times n}\hat{V}^T_{n\times n} = \begin{bmatrix} | & | & | & | \\ u_1 & u_2 & \dots & u_n \\ | & | & | & | \end{bmatrix} \begin{bmatrix} \sigma_1 & 0 & \dots & 0 \\ 0 & \sigma_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sigma_n \end{bmatrix} \begin{bmatrix} - & v_1^T & - \\ - & v_2^T & - \\ - & \vdots & - \\ - & v_n^T & - \end{bmatrix} \tag{2.24}$$

but $V^TV = I$, orthogonality allows:

$$A_{m\times n}\hat{V}_{n\times n} = \hat{U}_{m\times n}\hat{\Sigma}_{n\times n} \tag{2.25}$$

so,

$$A_{m\times n}\begin{bmatrix} | & | & | & | \\ v_1 & v_2 & \dots & v_n \\ | & | & | & | \end{bmatrix} = \begin{bmatrix} | & | & | & | \\ u_1 & u_2 & \dots & u_n \\ | & | & | & | \end{bmatrix} \begin{bmatrix} \sigma_1 & 0 & \dots & 0 \\ 0 & \sigma_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sigma_n \end{bmatrix} \tag{2.26}$$

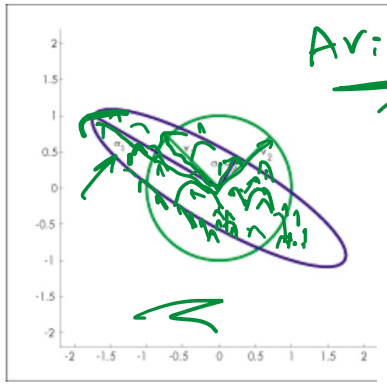
but this just states n-matrix times vector statements:

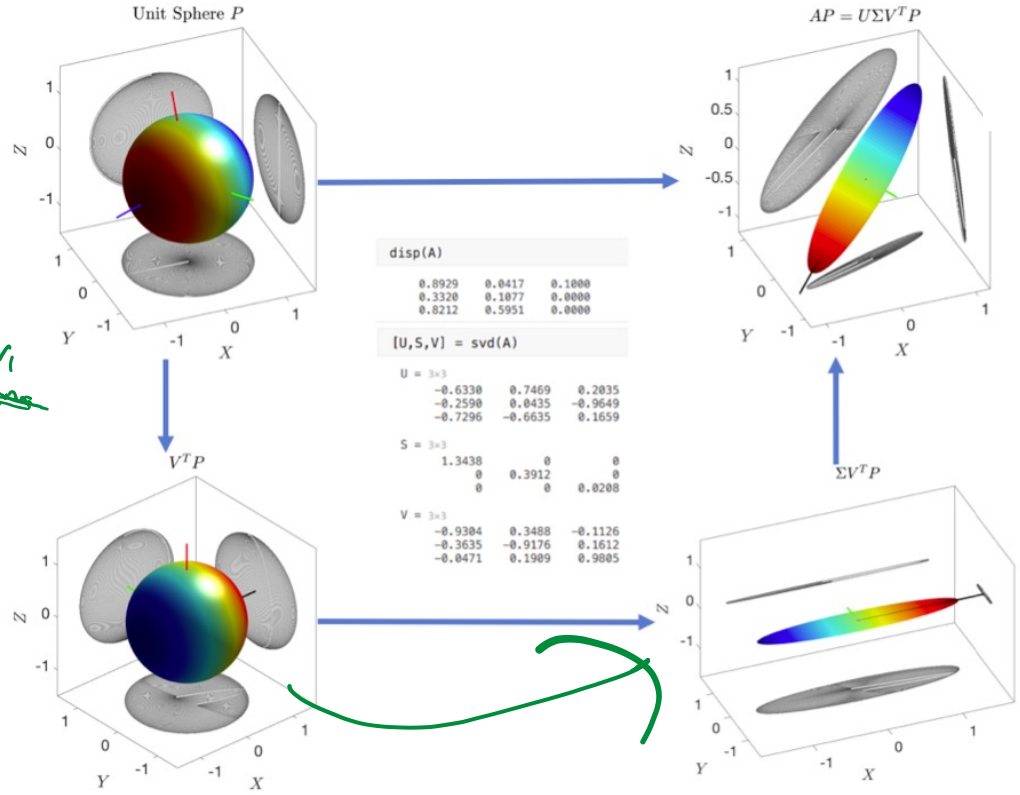
$$\begin{aligned} Av_1 &= \sigma_1 u_1 \\ Av_2 &= \sigma_2 u_2 \\ &\vdots \\ Av_n &= \sigma_n u_n \end{aligned} \tag{2.27}$$

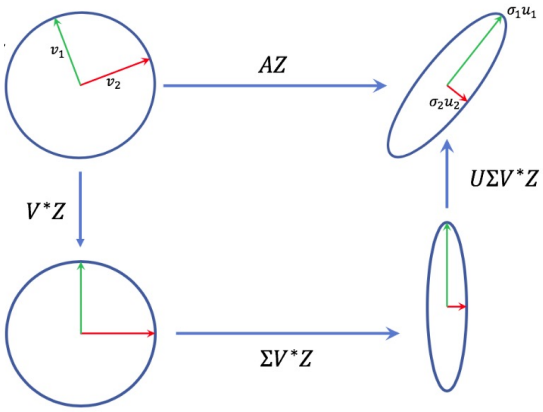
Full, Economy, Truncated SVD

Geometry:.

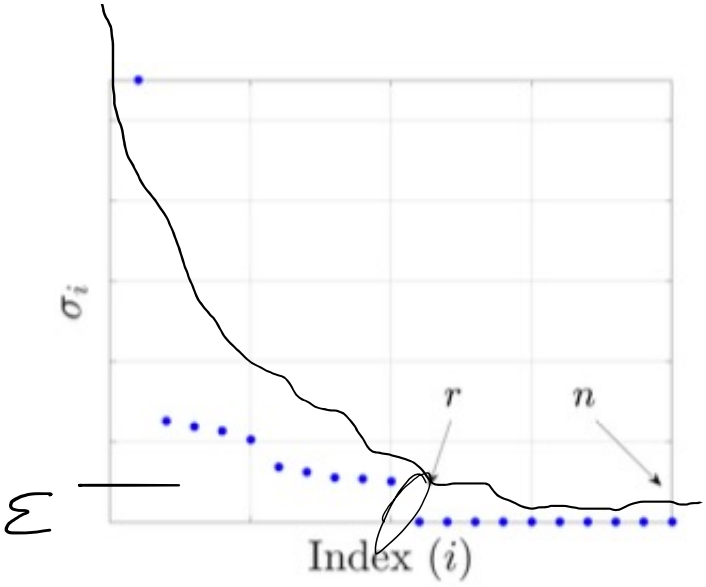
1. V^* rotates to a standard configuration.
2. Σ stretches each orthogonal axis to the major covariance axis of the corresponding ellipsoid, and
3. U rotates results back to the configuration that associates with A.







And ROM

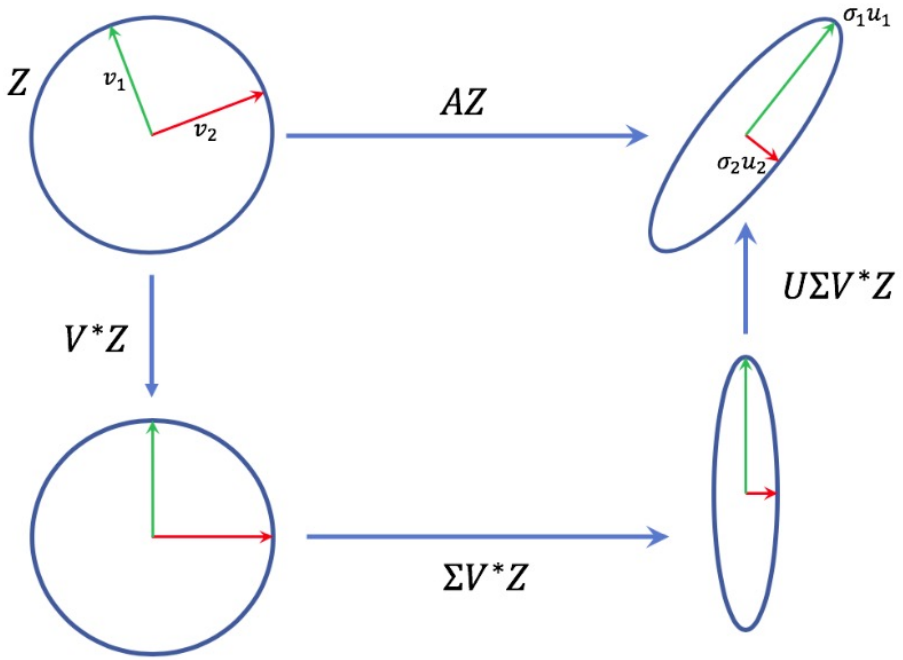


ε

$\sigma_r > \varepsilon > \sigma_{r+1}$

$\sim$

0



Bunny Compression

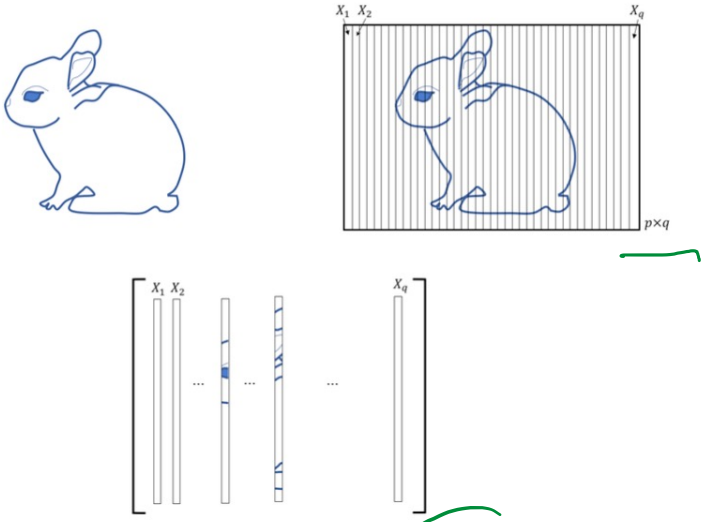
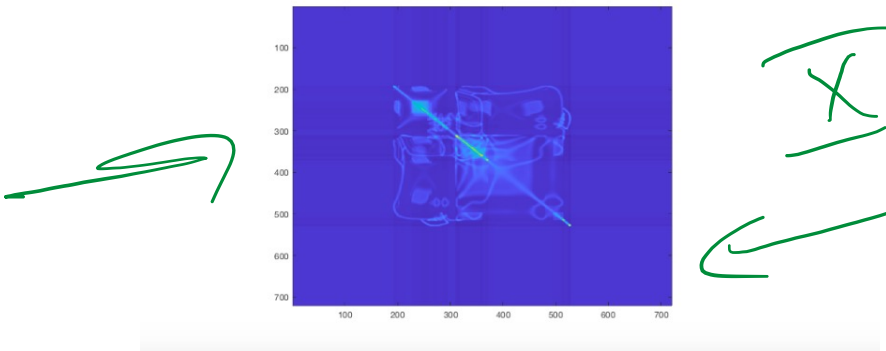
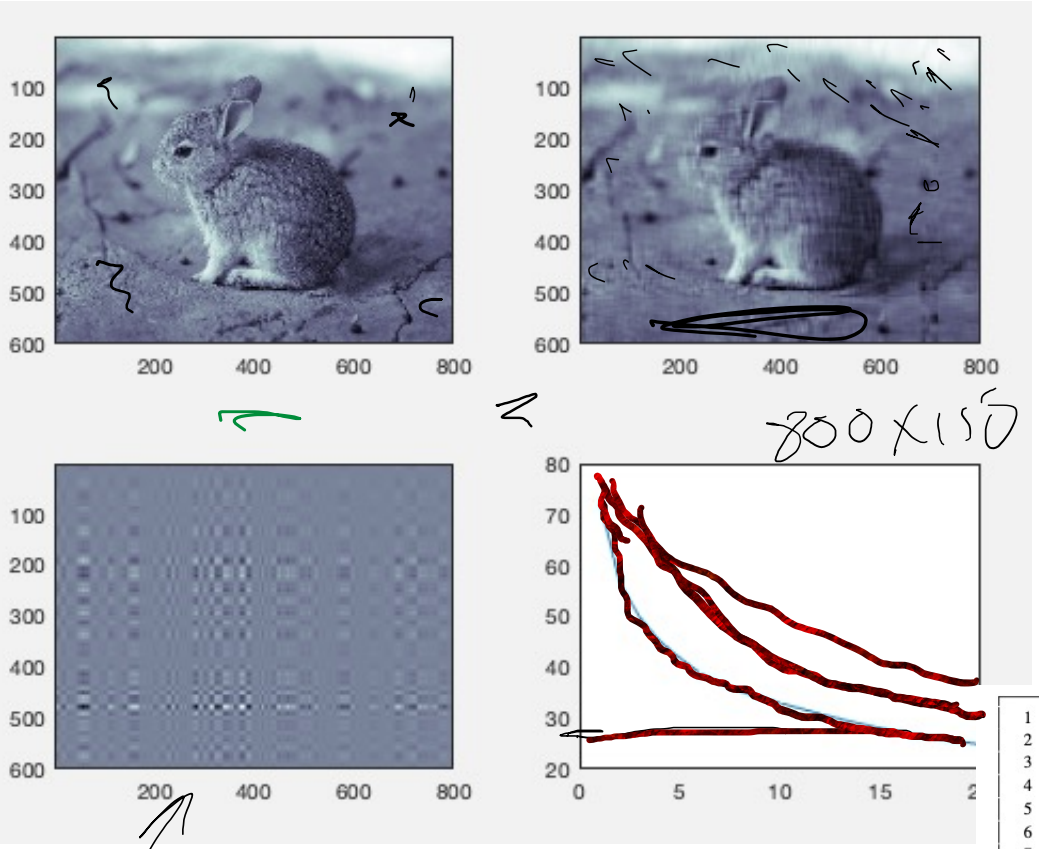


Figure 2.6: Caption

Covariance – notice the demean step

$$C_I = \frac{1}{n-1}\left(X - \tilde{X}\right)^T \left(X - \tilde{X}\right)$$





$$u(x,t) = \sum a_i(t)\phi_i(x)$$

$$|a_i(t)| \xrightarrow{i \to \infty} 0$$

as fast as possible.

```
I = imread('Bunny.jpg');

figure
subplot(1,2,1)
imshow(I)
xticks({}); yticks({});
pbaspect([1 1 1])
title('RGB Image')

I = rgb2gray(I);  %Convert the 3D RGB color to 1D grayscale
I = im2double(I); %Convert integer value to double (scaled ...
    from 0 to 1)

subplot(1,2,2)
imshow(I)
xticks({}); yticks({});
pbaspect([1 1 1])
title('Grayscale Image')
```

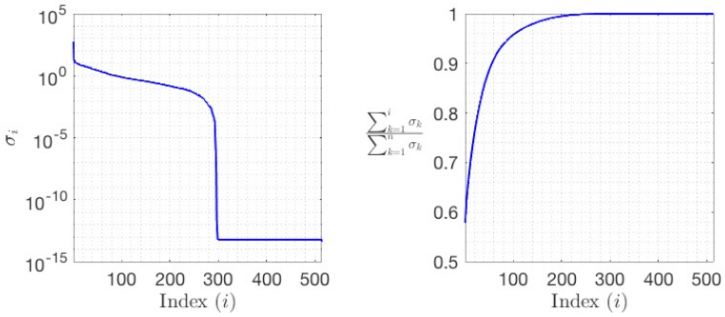
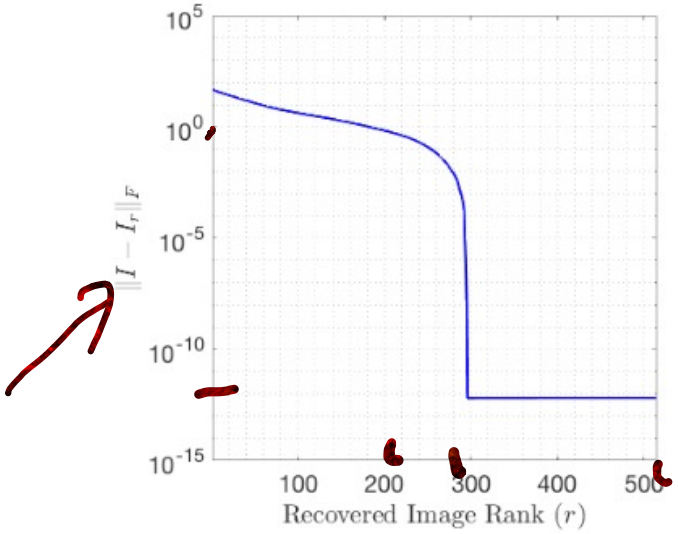
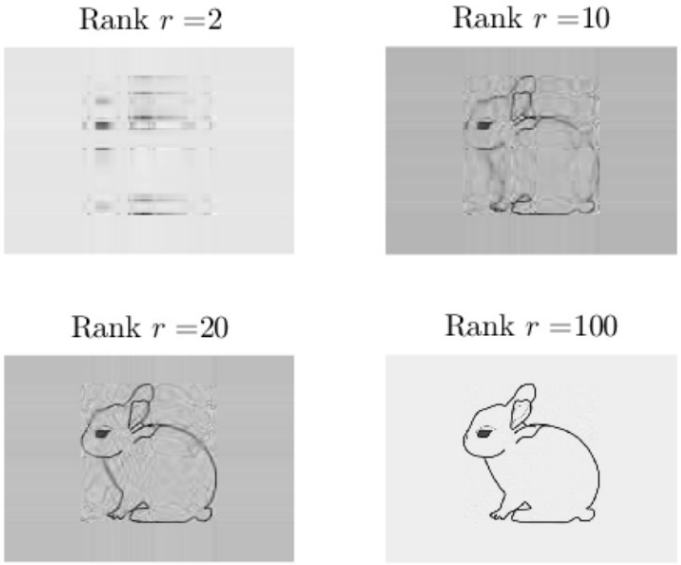


Figure 2.8: (Left) Singular Values. (Right) Energy





: Distance $\|I - I_r\|_F$, where I_r is the recovered image using the reduced

Code 2.1: Read, convert, and display images.

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```

History

Gene Golub's license plate, photographed by Professor P. M. Kroonenberg of Leiden University.Gene Howard Golub (February 29, 1932 – November 16, 2007), Fletcher Jones Professor of Computer Science at Stanford University. His work made fundamental contributions that have made the singular value decomposition practical as one of the most powerful and widely used tools in modern matrix computation.

Lots of Machine learning
& Data Analysis
is solving an ill-posed
- optimize a cost function.

$$AA^T = U \Sigma V^T (V \Sigma^T U^T)$$

$$= U \Sigma \Sigma^T U^T$$

$$(AA^T)\underline{U} = U(\Sigma\Sigma^T) = (\Sigma\Sigma^T)\underline{U}$$

$$\underline{U} = (u_1 \; u_2 \; \ldots \; u_m)$$

Definition 2.1.2 — Induced Norm. Suppose a vector norm $\|\cdot\|$ on $\mathcal{K}^m$ is given. Any matrix $A_{m\times n}$ induces a linear operator from $\mathcal{K}^n$ to $\mathcal{K}^m$ with respect to the standard basis, and one defines the corresponding induced norm or operator norm on the space $\mathcal{K}^{m\times n}$ of all $m \times n$ matrices as follows:

$$\|A\|_p = \sup_{x\neq 0} \frac{\|Ax\|_p}{\|x\|_p} \tag{2.14}$$

or, taking a vector x such that $\|x\|_p = 1$, then we have

$$\|A\|_p = \sup_{\|x\|_p=1} \|Ax\|_p \tag{2.15}$$

Some Special (Simple) Matrix Norms

The first 3 of these are induced norms, but the 4th is not.

- For $p = 1$:

$$\|A\|_1 = \max_{1\leq j\leq n} \sum_{i=1}^{m} |a_{ij}| \tag{2.16}$$

- For $p = \infty$:

$$\|A\|_\infty = \max_{1\leq i\leq m} \sum_{j=1}^{n} |a_{ij}| \tag{2.17}$$

- A special case is the spectral norm when $p = 2$, in which we have:

$$\|A\|_2 = \sqrt{\lambda_{max}(A^T A)} = \sigma_{max} \tag{2.18}$$

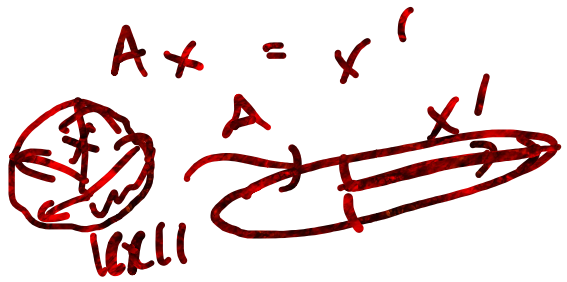
  where σ_{max} is the maximum singular value of the matrix A.

- The Frobenius norm is given by:

$$\|A\|_F = \sqrt{\sum_{i=1}^{m}\sum_{j=1}^{n} |a_{ij}|^2} = \sqrt{\sum_{i=1}^{\min\{m,n\}} \sigma_i^2} \tag{2.19}$$

Theorem 2.1.2 For a matrix A, the product of the singular values of A, equals the absolute value of its determinant:

$$|det(A)| = \prod_{i=1}^{n} \sigma_i \tag{2.20}$$



$p=1 : \|(x_1,x_2)\|_1 = |x_1| + |x_2|$

$p=\infty : \|(x_1,x_2)\|_\infty = \max_i |x_i|$

$p=2 : \|(x_1,x_2)\|_2 = \sqrt{x_1^2 + x_2^2}$

$x = \begin{pmatrix} 1 \\ 3 \end{pmatrix}$; $\|x\|_1 = 1+3 = 4$

$\|x\|_\infty = 3$

$\|x\|_2 = \sqrt{1^2+3^2} = \sqrt{10}$

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Fun facts about matrix estimation (data estimation)

If A, $\sigma_1 \geq \cdots \geq \sigma_r > \sigma_{r+1} = 0$

- $\text{range}(A) = \text{span}(u_1, u_2, \ldots, u_r)$
- $\text{null}(A) = \text{span}(v_{r+1}, v_{r+2}, \ldots, v_n)$

$0 \xrightarrow{A}$ ____ $\sigma_2 = 0$, $r = 1$

- $\|A\|_2 = \sigma_1$; $\|A\|_F = \sqrt{\sigma_1^2 + \sigma_2^2 + \cdots + \sigma_r^2}$
- $A = \sum_{i=1}^{r} \sigma_i u_i v_i^* = \sigma_1 u_1 v_1^* + \sigma_2 u_2 v_2^* + \cdots + \sigma_r v_r^*$

rank-1 outer products.

$w_1^T w_2 = w_1 \cdot w_2 = \|w_1\| \|w_2\| \cos\theta$

$A = U \Sigma V^T$

$\begin{pmatrix} | & | & & | \\ u_1 & u_2 & \cdots & u_m \\ | & | & & | \end{pmatrix}_{m \times r} \begin{pmatrix} \sigma_1 & & & \\ & \sigma_2 & & \\ & & \ddots & \\ & & & \sigma_r \end{pmatrix} \begin{pmatrix} - & v_1^T & - \\ - & v_2^T & - \\ & \vdots & \end{pmatrix}_{r \times n}$

$w_1 \to w_2$

$\begin{pmatrix} 1 & 2 \\ 2 & 4 \end{pmatrix}$

Matrix Estimation / Data Estimation. $A_{m \times n}$

A

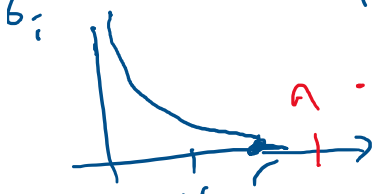
- let $\quad 0 \le N \le r$ and $A_N = \sum_{i=1}^{N} \sigma_i u_i v_i^*$

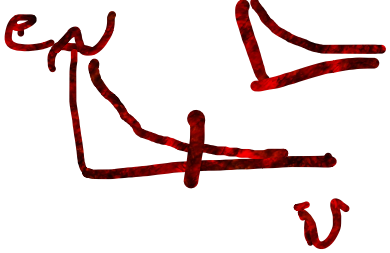
(so we may be skipping some of them ... $\sum_{i=N+1}^{r}$)

Then

- $\| A - A_N \|_2 = \sigma_{N+1}$ (first one skipped)

(what if it zero?)





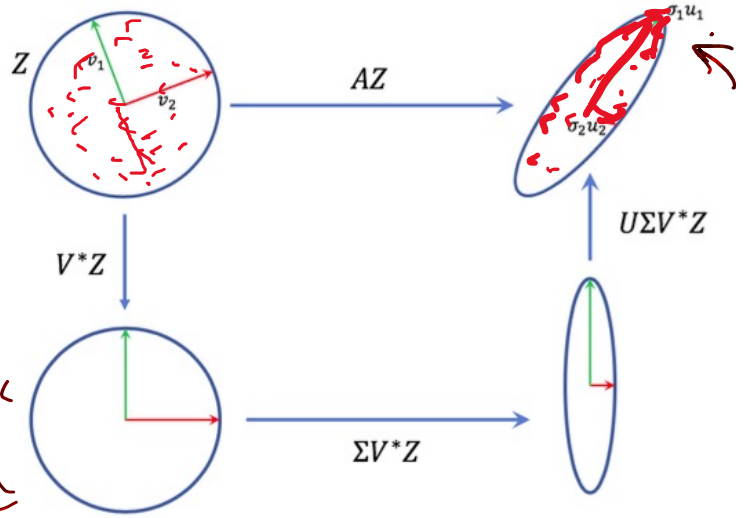
- $\| A - A_N \|_F = \sqrt{\sigma_{N+1}^2 + \sigma_{N+2}^2 + \cdots + \sigma_r^2}$

On PCA Principal Component Analysis, Eigenface

-On Raleigh Ritz Quotient

-On Spectral Decomposition Theorem

-On Data Clouds



$\bar{X} = [x_1 | x_2 | \dots | x_n]$ $\vec{x}_i = a_i \vec{v}_1 + a_2 \vec{v}_2 + \dots + a_r \vec{v}_r$

v_i's as basis set.

x_i = PCA gives basis set where $\langle a_i \rangle$ as fast as possible vs any other basis set.

$\langle a_i \rangle = \frac{1}{n} \sum_{i=1}^{n} a_i$

Data for PCA - "Pretend data looks like an ellipsoid"

Ex: $X_i \sim$ 4000 x 1 gene expression table for each i.

$i = 1 \ldots 216$ patients

$$X_i = \begin{pmatrix} x_i^1 \\ x_i^2 \\ \vdots \\ x_{4000} \end{pmatrix}$$

$y_i = 0$ or 1 "0" if not cancer "1" if cancer.

$f: \mathbb{R}^{4000} \to \mathbb{Z}_2$ $\quad \mathbb{Z}_2 = \{0, 1\}$.

$x \in \mathbb{R}^{4000}$

- Supervised vs. unsupervised.

unsupervised - just input - just structure/geometry.

"Data Cloud" ~ distribution of R.V. $x \sim X$

just x

$\rho(x): \mathbb{R}^{4000} \to \mathbb{R}^+$ a

- supervised learning is descriptive $f: x \to y$ (labels)

The spectral decomposition

Let A be a $n \times n$ symmetric matrix. From the spectral theorem, we know that there is an orthonormal basis $u_1, \cdots, u_n$ of $\mathbb{R}^n$ such that each u_j is an eigenvector of A. Let λ_j be the eigenvalue corresponding to u_j, that is,

$$Au_j = \lambda_j u_j.$$

Then

$$A = PDP^{-1} = PDP^T$$

where P is the orthogonal matrix $P = [u_1 \ \cdots \ u_n]$ and D is the diagonal matrix with diagonal entries $\lambda_1, \cdots, \lambda_n$. The equation $A = PDP^T$ can be rewritten as:

$$\begin{aligned} A &= [u_1 \ \cdots \ u_n] \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix} \begin{bmatrix} u_1^T \\ \vdots \\ u_n^T \end{bmatrix} \\ &= [\lambda_1 u_1 \ \cdots \ \lambda_n u_n] \begin{bmatrix} u_1^T \\ \vdots \\ u_n^T \end{bmatrix} \\ &= \lambda_1 u_1 u_1^T + \cdots + \lambda_n u_n u_n^T. \end{aligned}$$

The expression

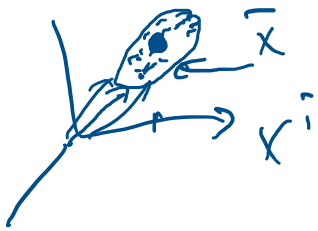
$$A = \lambda_1 u_1 u_1^T + \cdots + \lambda_n u_n u_n^T.$$

is called the spectral decomposition of A. Note that each matrix $u_j u_j^T$ has rank 1 and is the matrix of projection onto the one dimensional subspace spanned by u_j. In other words, the linear map P defined by $P(x) = u_j u_j^T x$ is the orthogonal projection onto the subspace spanned by u_j.

PCA as algorithm

(say n = 4000, n = 216)

- Data $\underline{X} = \begin{pmatrix} | & | & & | \\ x_1 & x_2 & \cdots & x_n \\ | & | & & | \end{pmatrix}_{m \times n}$ ← i^{th}



- what if $x_i \sim \mathcal{N}(\bar{x}, \Sigma)$ — covariance matrix.

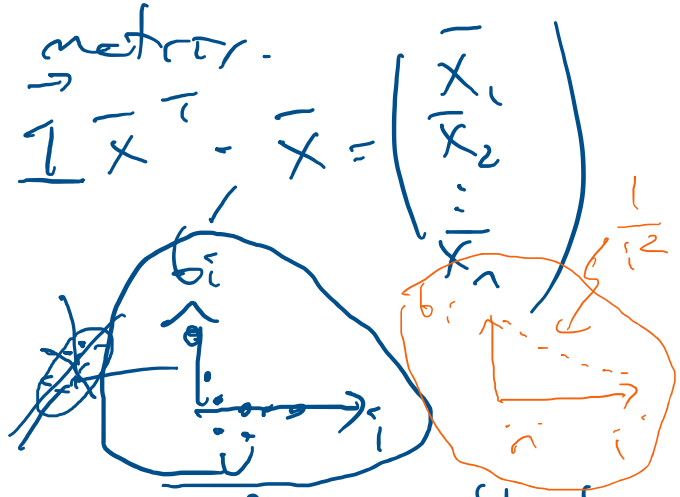
$$B = \underline{X} - \bar{B}; \quad \bar{B} = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}_{n \times 1} \bar{x}^T_{1 \times n} = \underline{1}\,\bar{x}^T \cdot \bar{x} = \begin{pmatrix} \bar{x}_1 \\ \bar{x}_2 \\ \vdots \\ \bar{x}_n \end{pmatrix}$$

$$\bar{x}_i = \frac{1}{n} \sum_{i=1}^{n} \underline{X}_{ij}$$

$$B = U \Sigma V^T; \quad U = [U_1, U_2 \cdots U_n]$$

and U_1 is major axis — most energetic feature that explains most data.

U_2 is first minor axis



Side note:

$\frac{1}{i^2}$ is slowest converging to zero σ_i^2

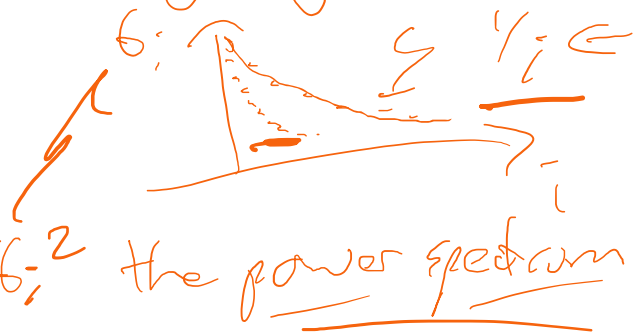
i.e. $\frac{1}{i}$ is slowest converging σ_i

$\sigma_i \le \frac{1}{i}$

$\sum_{i=1}^{\infty} \frac{1}{i^p} < \infty$ if bad.

$p < 1$, $p = 1$ harmonic

$= 1 + \frac{1}{2^p} + \frac{1}{3^p} + \frac{1}{4^p} + \cdots \quad p = 1$



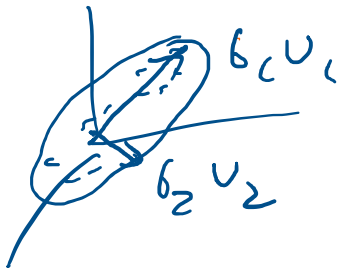
the power spectrum

$B = I\bar{X}^T$; let $C = \frac{1}{n-1} B^T B$ covariance

it covariance matrix

$B = U \Sigma V^T$

$U = [u_1 | u_2 | \dots | u_n]$



$\sigma_3 \approx 0$

$\sigma_2 \Sigma 0$

- $u_1 = \underset{\|u\|=1}{\operatorname{argmax}} \; u^T B^T B u = \underset{u}{\operatorname{argmax}} \frac{u^T B^T B u}{u^T u}$

  $\|Bu\|_2^2$

  Rayleigh-Ritz quotient - $Bu \cdot Bu = \|Bu\|_2^2$

- $u_2 = \underset{\substack{\|u\|=1 \\ u \perp u_1}}{\operatorname{argmax}} \; u^T B^T B u$

$$
\begin{aligned}
||x_i - proj_w x_i||_2^2 &= ||x_i(w \cdot x_i)w||_2^2 = (x_i - (w^T x_i)w)^T(x_i - (w^T x_i)w) \\
&= (x^T x_i) - (w^T x_i)^2 = ||x_i||_2^2 - (w^T x_i)^2. \qquad (2.43)
\end{aligned}
$$

To minimize this residual with respect to the unknown vector w, averaged across the data set, it is sufficient to maximize the second term since the first term does not depend on w. Thus we wish to maximize,

$$
\mathcal{L}_1(\mathcal{D}; \Theta) = \frac{1}{n} \sum_{i=1}^{N} (w^T x_i)^2, \qquad (2.44)
$$

The Eigs of C=B'B give optimal projection – thus PCA and…. KL

- $CV = VD$ eigenvectors of C all stacked.

Ax optimize $x^T Ax$, maximize.

$x = \begin{pmatrix} x^1 \\ \vdots \\ x^j \\ \vdots \\ x^n \end{pmatrix}$ $r(x) = \frac{x^T Ax}{x^T x}$ ($A = B^T B$, $x = 0$)

$$\frac{\partial r}{\partial x^j} = \frac{\partial}{\partial x^j}\left(r(x)\right) = \frac{\frac{\partial}{\partial x^j}(x^T Ax) - x^T Ax \frac{\partial}{\partial x^j}(x^T x)}{(x^T x)^2}$$

$$= \frac{2(Ax)_j}{x^T x} - \frac{(x^T Ax)\,2x_j}{(x^T x)^2}$$

$$\nabla r(x) = \frac{2}{x^T x}\left(Ax - r(x)x\right) = \frac{2}{x^T x}\left(A - r(x)\right)x = 0$$

$Ax = r(x)\,x \Rightarrow Ax = \lambda x$

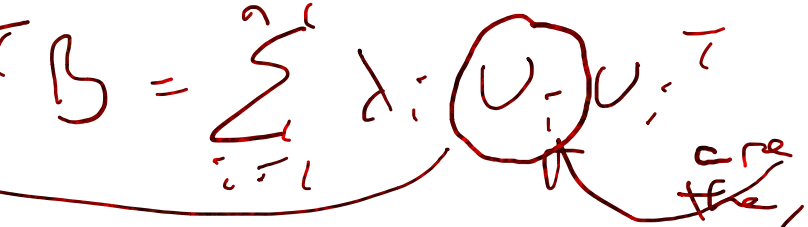
$r(x)$ = λ

Conclude

The x that optimizes $r(x) = \frac{x^T A x}{x^T x}$

is an eigenvector and $r(x)$ is its eigenvalue.

- S.D.T. for $A = B^T B = \sum_{i=1}^{r} \lambda_i v_i v_i^T$



let $B = U \Sigma V^T$ (svd)

$B^T B = V \Sigma^T U^T U \Sigma V^T = V \Sigma^T \Sigma V^T$

$D = \begin{pmatrix} \sigma_1 & & & \\ & \ddots & & \\ & & \sigma_r & \\ & & & 0 \end{pmatrix}$ $\begin{pmatrix} \sigma_1 & & & \\ & \ddots & & \\ & & \sigma_r & \\ & & & 0 \end{pmatrix} = \begin{pmatrix} \sigma_1^2 & & & \\ & \sigma_2^2 & & \\ & & \ddots & \\ & & & \sigma_r^2 \\ & & & & 0 \end{pmatrix}$

$= V D V^T = \sum_{i=1}^{r} \sigma_i^2 v_i v_i^T$

$$\mathcal{L}_1(\mathcal{D};\Theta) = \frac{1}{n}(X^T w)^T(X^T w) = \frac{1}{n} w^T(XX^T)w, \tag{2.45}$$

and the matrix $\frac{1}{n}(XX^T)$ is familiar in statistics as a covariance matrix. To optimize $\mathcal{L}_1$, subject to a constraint,[12]

$$\|w\|_2 = 1, \tag{2.46}$$

we can use the Lagrange multiplier method by defining an expanded loss function(cost function) with the equality constraint built in with a Lagrange multiplier. Let,

$$\mathcal{L}(\mathcal{D};\Theta,\lambda) = \mathcal{L}_1(\mathcal{D};\Theta) - \lambda(w^T w - 1) = \frac{1}{n} w^T(XX^T)w - \lambda(w^T w - 1). \tag{2.47}$$

To minimize this, we take derivatives and set them equal to zero.

$$\frac{\partial \mathcal{L}}{\partial w} = \frac{2}{n} X X^T w - 2\lambda w \implies \frac{1}{n}(X X^T) w = \lambda w \tag{2.48}$$

$$\frac{\partial \mathcal{L}}{\partial \lambda} = w^T w - 1 \implies \|w\|_2 = 1. \tag{2.49}$$

Theorem 2.2.1 — PCA foundations. Let A be a symmetric $d \times d$ matrix. Then its (real) eigenvalues $\lambda_1 \geq \lambda_2 \geq ... \geq \lambda_d$, associate with orthogonal eigenvectors $w_1, w_2, ..., w_d$. Furthermore,

$$\begin{aligned}
\lambda_1 &= \max_{\|w\|=1} w^T A w, \text{ with } w_1 = \arg\max_{\|w\|=1} w^T A w.r \\
\lambda_2 &= \max_{\|w\|=1, w \perp w_1} w^T A w \text{ with } w_2 = \arg\max_{\|w\|=1, w \perp w_1} w^T A w. \\
&\vdots \\
\lambda_d &= \max_{\|w\|=1, w \perp w_1, w_2, .., w_{d-1}} w^T A w \text{ with } w_d = \arg\max_{\|w\|=1, w \perp w_1, w_2, .., w_{d-1}} w^T A w.
\end{aligned} \tag{2.50}$$

Theorem 2.2.2 — Spectral Decomposition. If A is a symmetric positive semi-definite matrix, then there is an orthogonal set of eigenvectors, u_i, each with non-negative eigenvalues, $\lambda_i \geq 0$. Furthermore, the decomposition of A has the following representation by rank one matrices that describe the action of A as a weighted sum of simple projections onto the subspaces spanned by each u_i,

$$A = \sum_{i=1}^{N} \lambda_i u_i u_i^T \tag{2.51}$$

Which functions are most efficient?

That is, we write a linear combination,

$$u(x,t) = \sum_k a_k(t)\varphi_k(x), \tag{2.57}$$

of functions $\varphi_k(x)$, where the time varying (component projection) values,

$$a_k(t) = \frac{(u(x,t), \varphi_k(x)}{\|\varphi_k(x)\|_2}, \tag{2.58}$$

or better yet, we can effectively skip the denominator by choosing the basis set of functions such that,

$$\|\varphi_k(x)\|_2^2 = (\varphi_k(x), \varphi_k(x)) = 1. \tag{2.59}$$

Given a spatiotemporal data sample as an array, (for example, typically from a solution derived from a computational solver for a PDE):

$$\mathbf{U} = \begin{pmatrix} | & | & & | \\ u(\vec{x}, t_1) & u(\vec{x}, t_2) & \dots & u(\vec{x}, t_T) \\ | & | & & | \end{pmatrix}. \tag{2.61}$$

then develop the demeaned array $U - \overline{U}$. **Find**:

$$\mathbf{\Phi} = \begin{pmatrix} | & | & & | \\ \varphi_1(x) & \varphi_2(x) & \dots & \varphi_k(x) \\ | & | & & | \end{pmatrix}. \tag{2.62}$$

by the singular value decomposition. This basis yields the fastest decaying power spectrum, *in time average*, versus all other possible basis.

$$
\begin{aligned}
a(t) &= \frac{(u,\varphi)}{\|\varphi\|^2} = \frac{\|u\|\|\varphi\|\cos\theta}{\|\varphi\|^2} = \|u\|\cos\theta \\
&= \int_\Omega u(x,t)\varphi(x)dx
\end{aligned} \tag{2.63}
$$

when $\|\varphi\| = 1$.
So, we will also write time average using brackets $< \cdot, \cdot >$, so define:

$$
\begin{aligned}
< |a(t)| > &= \frac{1}{T}\int_0^T |a(t)|dt \\
&= \frac{1}{T}\int_0^T |(u,\varphi)|dt \\
&= \frac{1}{T}\int_0^T \left|\int_\Omega u(x,t)\varphi(x)\right| dt
\end{aligned} \tag{2.64}
$$

$\star\star$ The goal is to choose a basis with fastest decaying power spectrum, in time average $\star\star$

Our goal can be summarized by the following loss function.

$$
\mathcal{L}(\varphi) = \frac{< |(u,\varphi)|^2 >}{\|\varphi\|^2}. \tag{2.65}
$$

This very compact notation, encodes two integrations. We remind that the round brackets describe the inner product, (f, g), meaning integration in the “space" variable. Now we have introduced the pointy brackets to describe time average. So,

$$
\mathcal{L}(\varphi) = \frac{\frac{1}{T}\int_0^T \left|\int_\Omega u(x,t)\varphi(x)dx\right|^2 dt}{\int_0^L \varphi^2(x)dx} \equiv < a\varphi^2(t) > \tag{2.66}
$$

or

$$
\max_{\|\varphi\|=1} \frac{1}{T}\int_0^T \left|\int_0^L u(x,t)\varphi(x)dx\right|^2 dt \tag{2.67}
$$

Theorem 2.3.1 — Parseval's Like Idenfity. If $f \in L^2([0, L])$, then if $\{\varphi_k(x)\}$ is an orthonormal basis set, then:

$$\|f\|_2^2 \leq \sum_{k=0}^{\infty} |a_k|^2 \tag{2.68}$$

where $a_k = f, \varphi_k)$.

$$\mathcal{L}(\varphi) =< |(u,\varphi)|^2 > -\lambda\left(\|\varphi\|^2 - 1\right)$$

$$C\vec{\varphi}(x) = \lambda\vec{\varphi}(x), C = \mathbf{U}\mathbf{U}^T,$$

$\frac{\partial}{\partial\delta}\mathcal{L}(\varphi + \delta\psi) = 0$:

$$\max_{\begin{cases}\|\varphi_1\| = 1\\ \varphi \in \mathcal{H}\end{cases}} < |(u,\varphi| >$$

and $\varphi_2(x)$ solves:

$$\max_{\begin{cases}\|\varphi_2\| = 1\\ \varphi \in \mathcal{H}\\ \varphi_2 \perp \varphi_1\end{cases}} < |(u,\varphi| >$$

Theorem 2.3.2 — Spectral Decomposition. If A is a symmetric positive semi-definite matrix (i.e. $\forall u \in \mathbb{R}^n, u^T A u \geq 0$), then there is an orthogonal set of (column) eigenvectors, v_i, each with non-negative eigenvalues, $\lambda_i \geq 0$, and furthermore the decomposition of A as the following rank one matrices describes the action of A as a weighted sum of simple projections onto the subspaces spanned by each v_i,

$$A = \sum_{i=1}^{N} \lambda_i v_i v_i^T \tag{2.71}$$

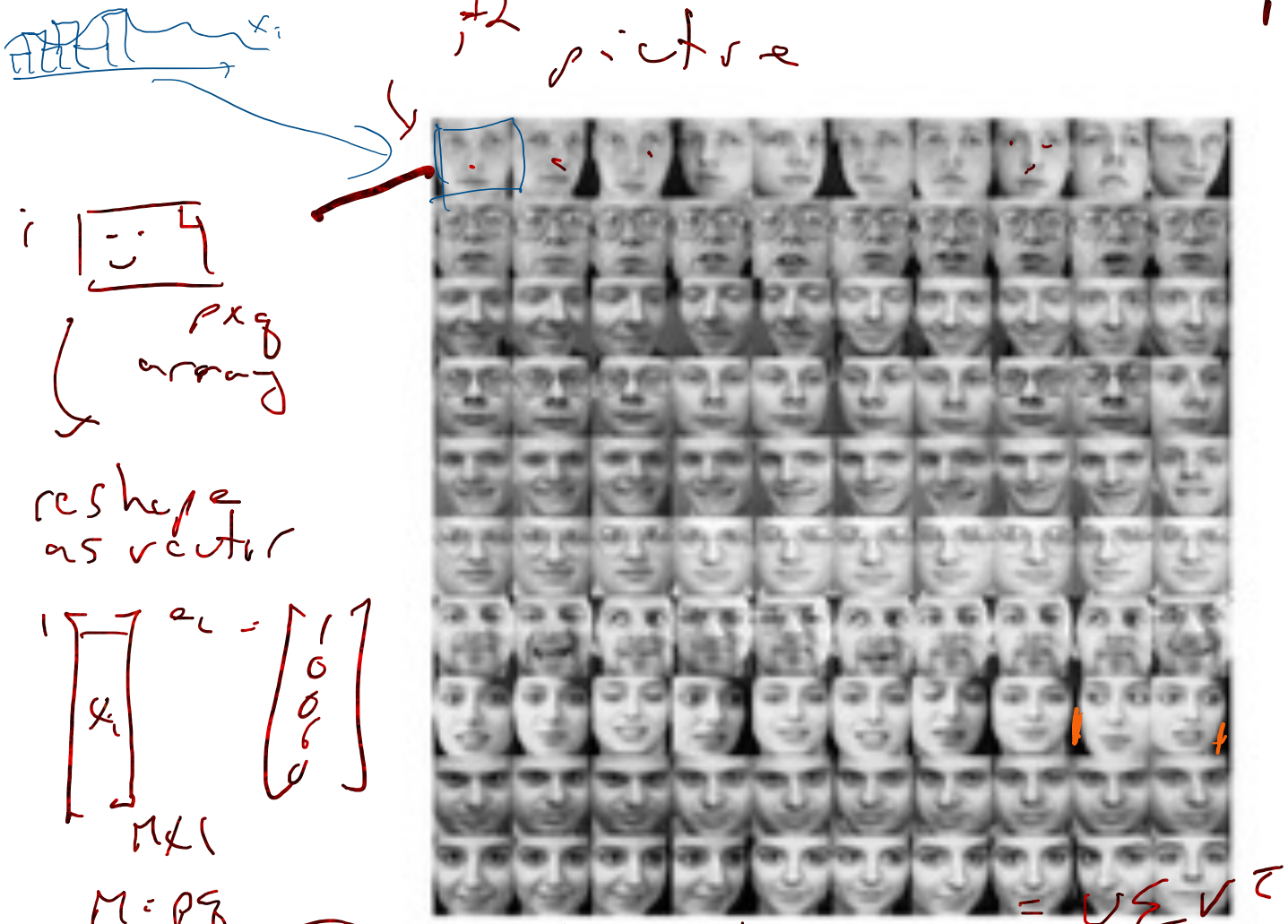
Solve $\mathbf{U}\mathbf{U}^t$ for eigs; solve fastest decaying time averaged power spectrum

$$\mathbf{U}\mathbf{U}^t V^t = \Lambda V^t, \tag{2.74}$$

we use $\mathbf{U} = U\Sigma V^t$, with $\Lambda = \Sigma^2$

Eigenface

08/31/20

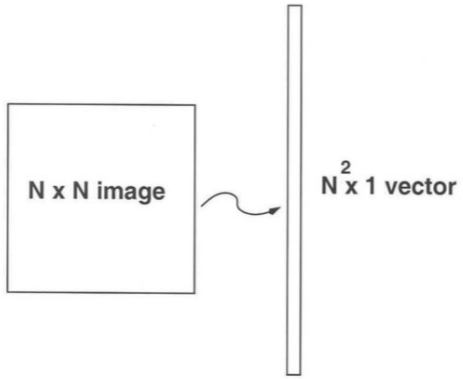


Eigenfaces for Face Detection/Recognition

(M. Turk and A. Pentland, "Eigenfaces for Recognition", *Journal of Cognitive Neuroscience*, vol. 3, no. 1, pp. 71-86, 1991, hard copy)

• Face Recognition

- The simplest approach is to think of it as a template matching problem:



- Problems arise when performing recognition in a high-dimensional space.

- Significant improvements can be achieved by first mapping the data into a *lower-dimensionality* space.

- How to find this lower-dimensional space?

• Main idea behind eigenfaces

- Suppose Γ is an N^2x1 vector, corresponding to an NxN face image I.

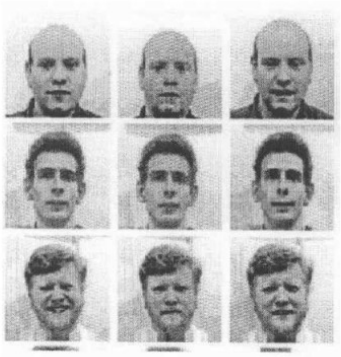
- The idea is to represent Γ (Φ=Γ - mean face) into a low-dimensional space:

$$\hat{\Phi} - mean = w_1 u_1 + w_2 u_2 + \cdots w_K u_K \ (K << N^2)$$

Computation of the eigenfaces

Step 1: obtain face images $I_1, I_2, ..., I_M$ (training faces)

(**very important:** the face images must be *centered* and of the same *size*)



Step 2: represent every image I_i as a vector Γ_i

Step 3: compute the average face vector Ψ:

$$\Psi = \frac{1}{M}\sum_{i=1}^{M}\Gamma_i$$

Step 4: subtract the mean face:

$$\Phi_i = \Gamma_i - \Psi$$

Step 5: compute the covariance matrix C:

$$C = \frac{1}{M}\sum_{n=1}^{M}\Phi_n\Phi_n^T = AA^T \quad (N^2 \text{x} N^2 \text{ matrix})$$

where $A = [\Phi_1\ \Phi_2 \cdots \Phi_M]$ $(N^2 \text{x} M \text{ matrix})$

Step 6: compute the eigenvectors u_i of AA^T

The matrix AA^T is very large --> not practical !!

Step 6.1: consider the matrix $A^T A$ (MxM matrix)

Step 6.2: compute the eigenvectors v_i of $A^T A$

$$A^T A v_i = \mu_i v_i$$

What is the relationship between us_i and v_i?

$A^T A v_i = \mu_i v_i \Rightarrow AA^T A v_i = \mu_i A v_i \Rightarrow$

$CAv_i = \mu_i Av_i$ or $Cu_i = \mu_i u_i$ where $u_i = Av_i$

Thus, AA^T and $A^T A$ have the same eigenvalues and their eigenvectors are related as follows: $u_i = Av_i$!!

Note 1: AA^T can have up to N^2 eigenvalues and eigenvectors.

Note 2: $A^T A$ can have up to M eigenvalues and eigenvectors.

Note 3: The M eigenvalues of $A^T A$ (along with their corresponding eigenvectors) correspond to the M *largest* eigenvalues of AA^T (along with their corresponding eigenvectors).

Step 6.3: compute the M best eigenvectors of AA^T: $u_i = Av_i$

(**important:** normalize u_i such that $\|u_i\| = 1$)

Step 7: keep only K eigenvectors (corresponding to the K largest eigenvalues)

Representing faces onto this basis

- Each face (minus the mean) Φ_i in the training set can be represented as a linear combination of the best K eigenvectors:

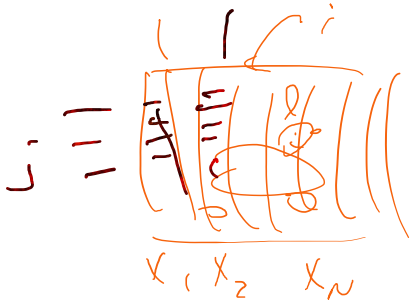
$$\hat{\Phi}_i - mean = \sum_{j=1}^{K} w_j u_j, \quad (w_j = u_j^T \Phi_i)$$

(we call the u_j's *eigenfaces*)

- Each normalized training face Φ_i is represented in this basis by a vector:

$$\Omega_i = \begin{bmatrix} w_1^i \\ w_2^i \\ \dots \\ w_K^i \end{bmatrix}, \quad i = 1, 2, \dots, M$$

Bunny



$\rightarrow \bar{X} = [x_1 \mid x_2 \mid \dots \mid x_N]$ ← one column

$= U \Sigma V^T$

PCA of bunny

$x_i = a_i^1 U_1 + a_i^2 U_2 + \dots + a_i^N U_N$

losses fall off as possible for the x_i set

$a_i = x_i^T U_i = \text{comp}_{U_i} x_i = \frac{U_i \cdot x_i}{|U_i|} = |x_i| \cos\theta$

On basis, functions, and Hilbert space.
Fourier, Taylor, Wavelet, POD-KL

leadin

3 B (2) – in 5 min. Signals analysis, Harmonic.

- Historically favorite basis set. $B = \{\omega_1, \omega_2, \ldots\}$

(vs. energy favorite basis set comes from PCA)

- Taylor polynomials. $f(x) = a_0 + a_1 x + a_2 x^2 + \ldots + a_r x^r$

$B = \{x^0, x^1, x^2, \ldots\}$

$a_i = \frac{f^{(i)}(0)(x-0)}{i!}$

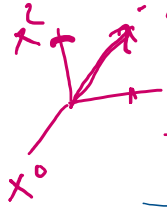
- Fourier modes.

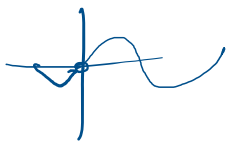
$f(x) = a_1 \sin x + a_2 \sin 2x + a_3 \sin 3x + \ldots$

$B = \{\sin x, \sin 2x, \ldots\}$

$a_i = (f(x), \sin x) = \frac{1}{2\pi} \int_0 f(x) \sin x \, dx$

$\mathcal{H}$: Hilbert space



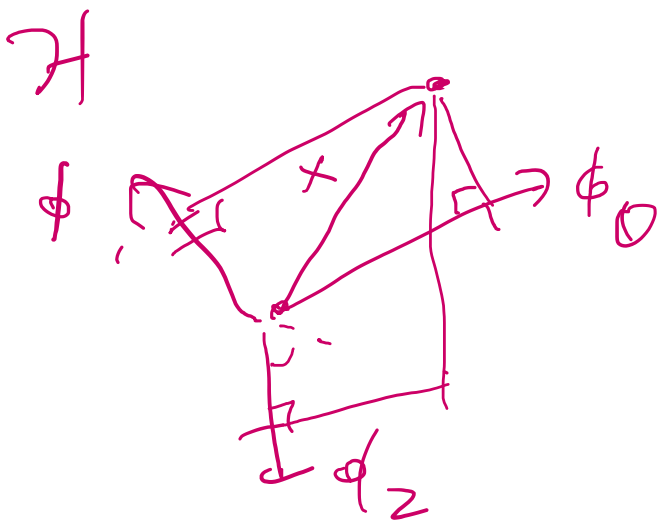


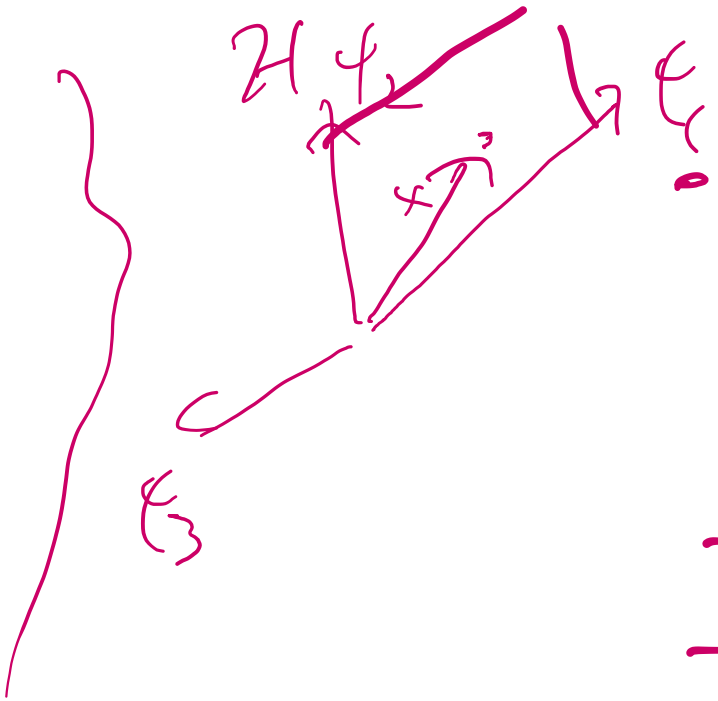
- Wavelet basis.
- chebychev polys, legendre ---

$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} + \ldots$

$\lim_{x \to 0} \frac{\sin x}{x} = 1$

Changing basis is a sort of coord-rot.





$\phi_0(y) = y^0$

$\phi_1(y) = y^1$

$\phi_2(y) = y^2$

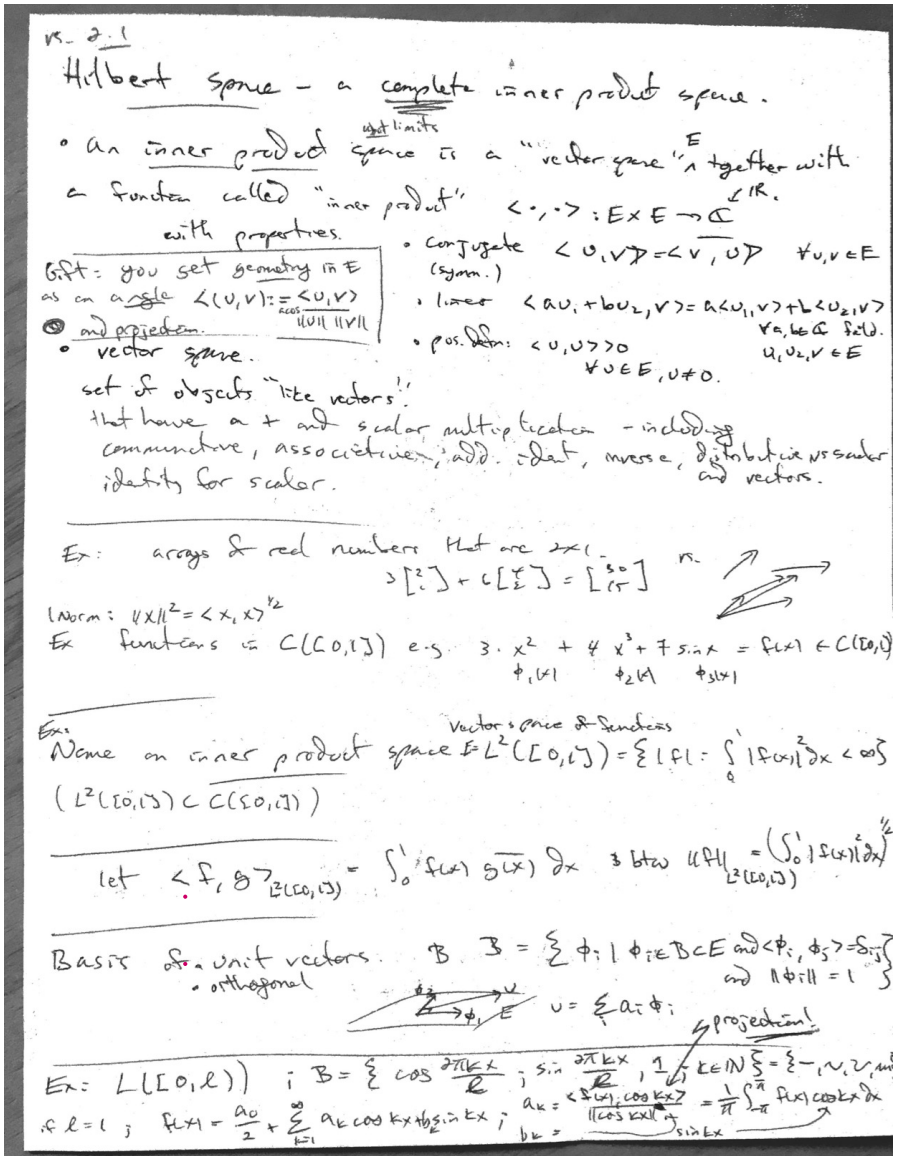
$\vdots$

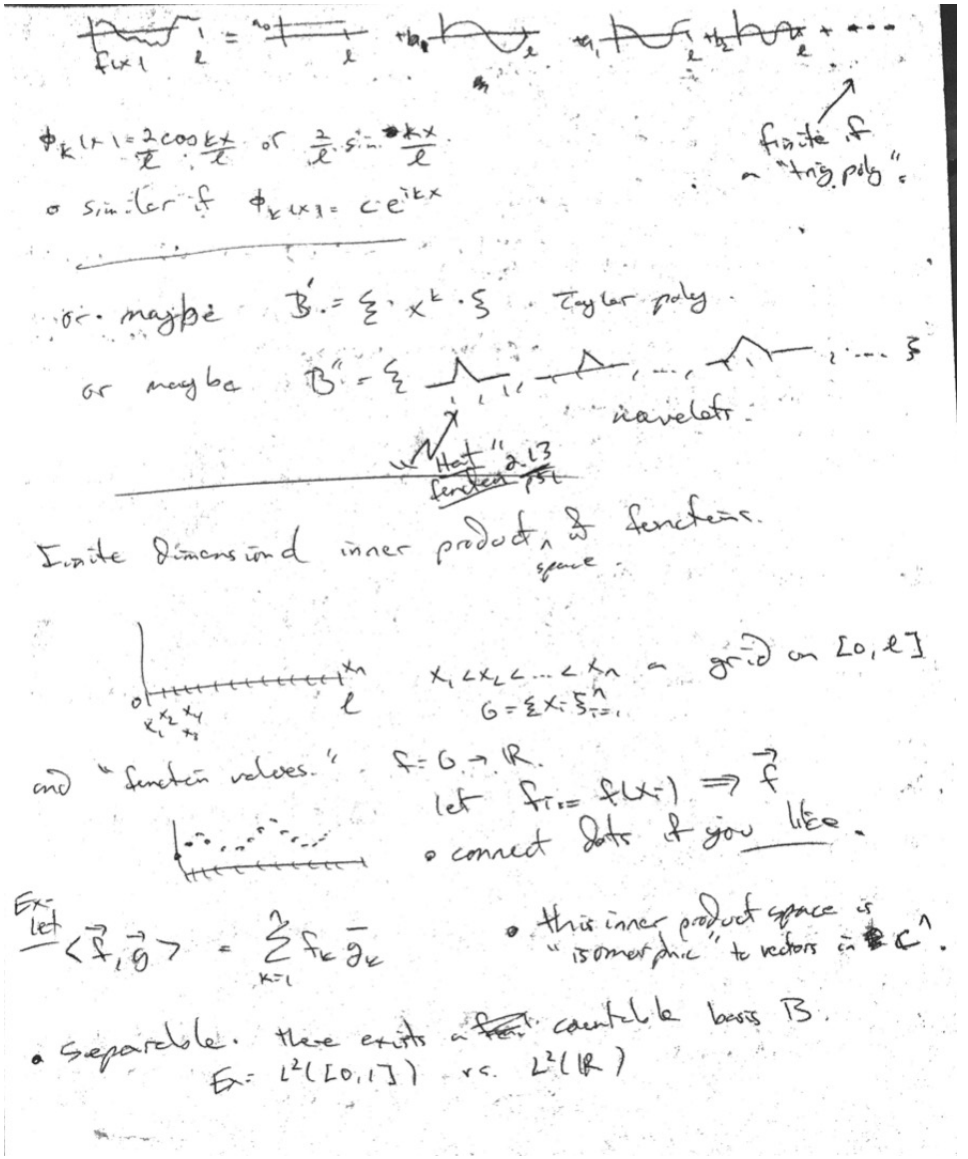
$\psi_1(y) = \sin(y)$

$\psi_2(y) = \sin(2y)$

$\psi_3(y) = \sin(3y)$

On basis, functions, and Hilbert space. Fourier, Taylor, Wavelet, POD-KL





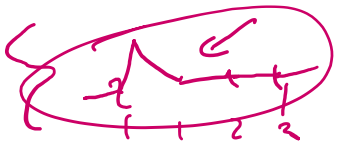
• $B = \{$  $\}$ [0,1] Fourier ⇒

Fourier $\{1, \sin x, \cos x, \sin 2x, \cos 2x, \cdots\}$

• $B' = \{$  $\} = \{1, x, x^2, x^3 \cdots\}$

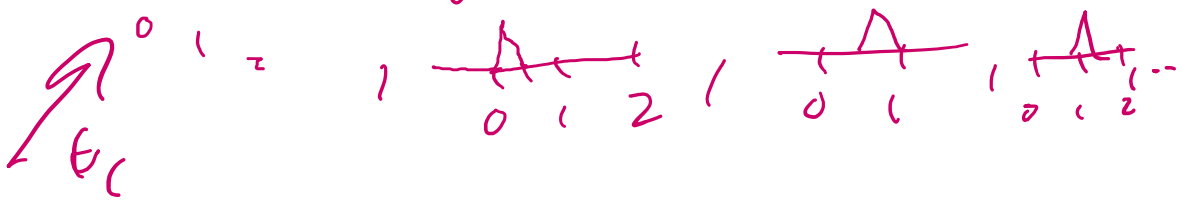
Taylor, ... Legendre,

Taylor.

• B'' = 

$\phi_2(x)$ $\phi_3(x)$

Wavelet.



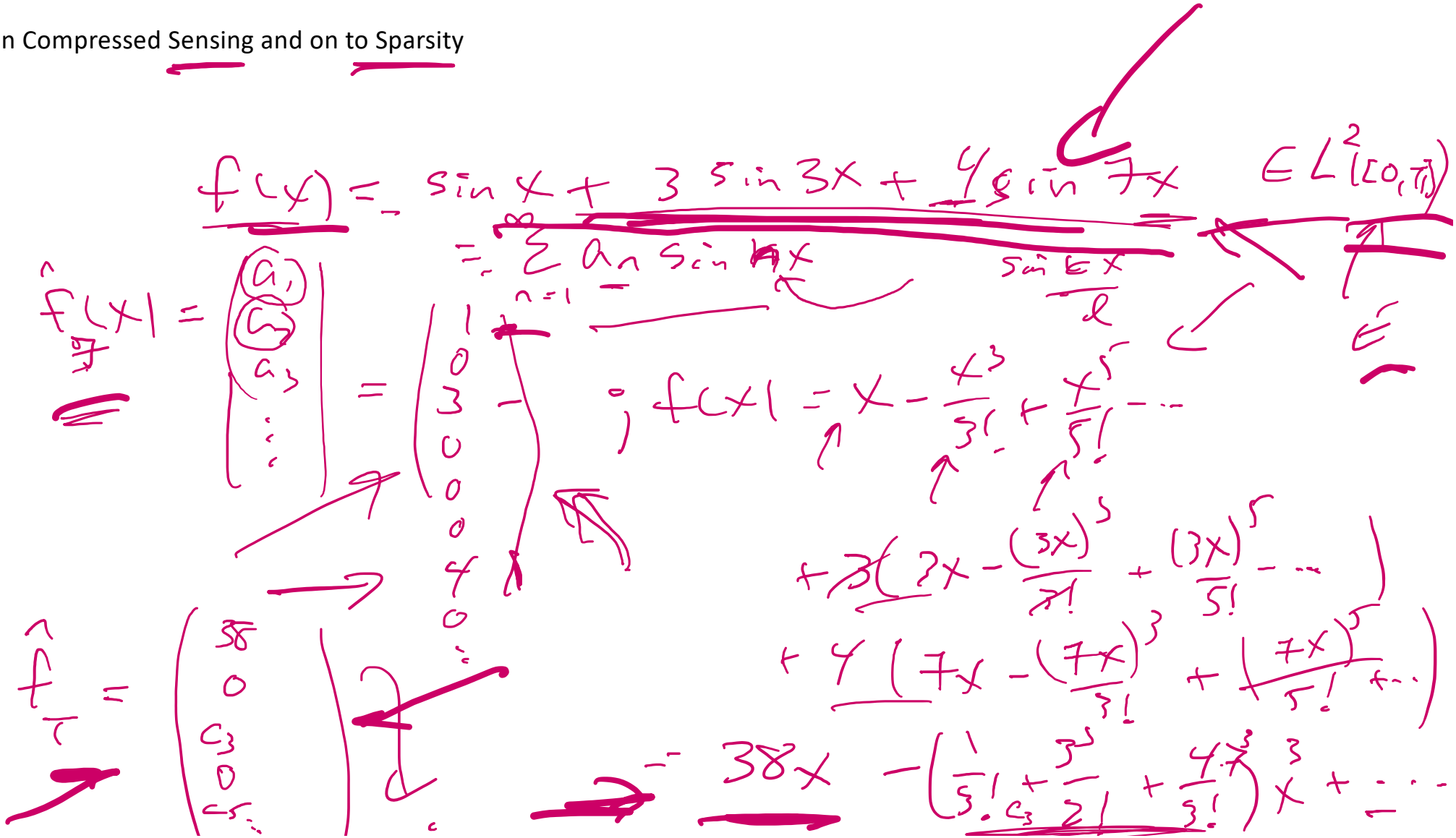
$\phi(x)$

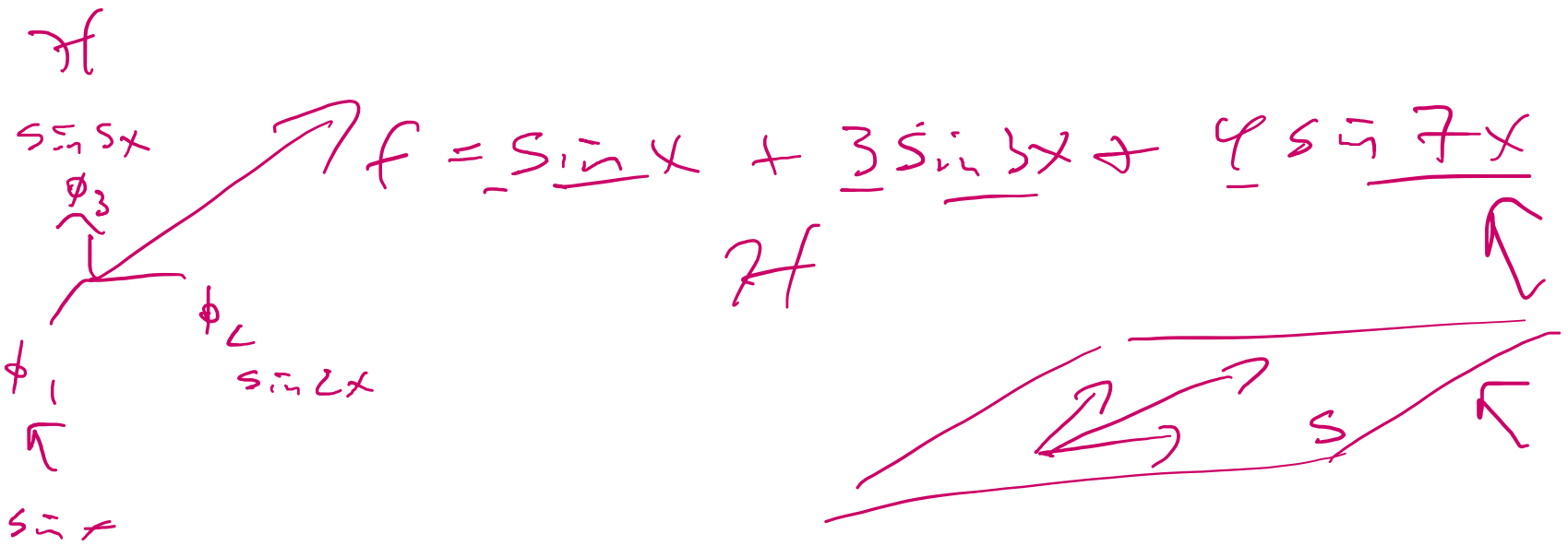
$f(x) \in L^2(\mathbb{R})$,

$f(x) = \sum a_i \phi_i(x)$; $a_i = \langle f, \phi_i(x) \rangle / \|\phi_i(x)\|$

•  Haar....

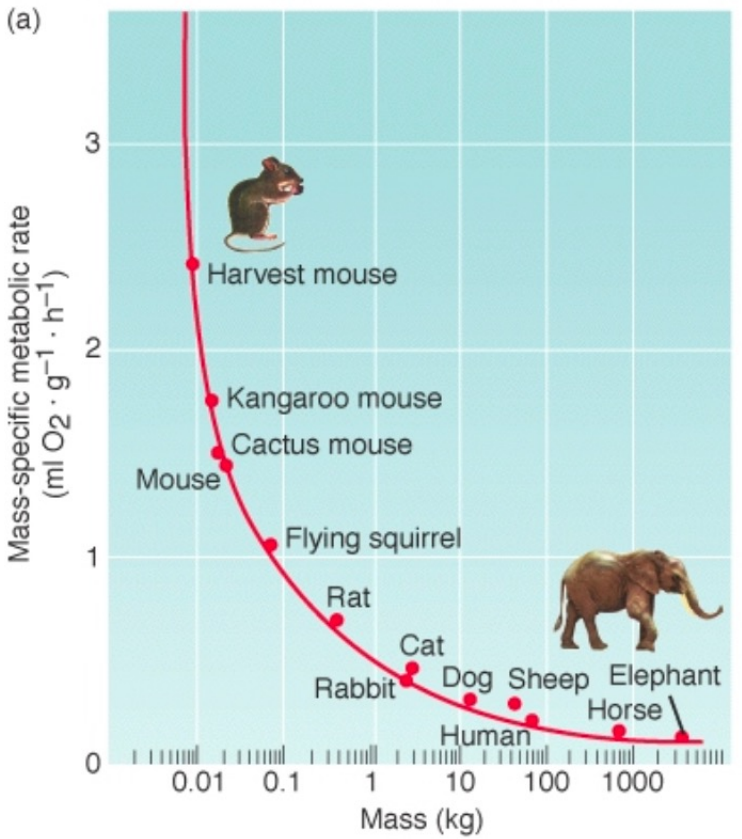
On Compressed Sensing and on to Sparsity





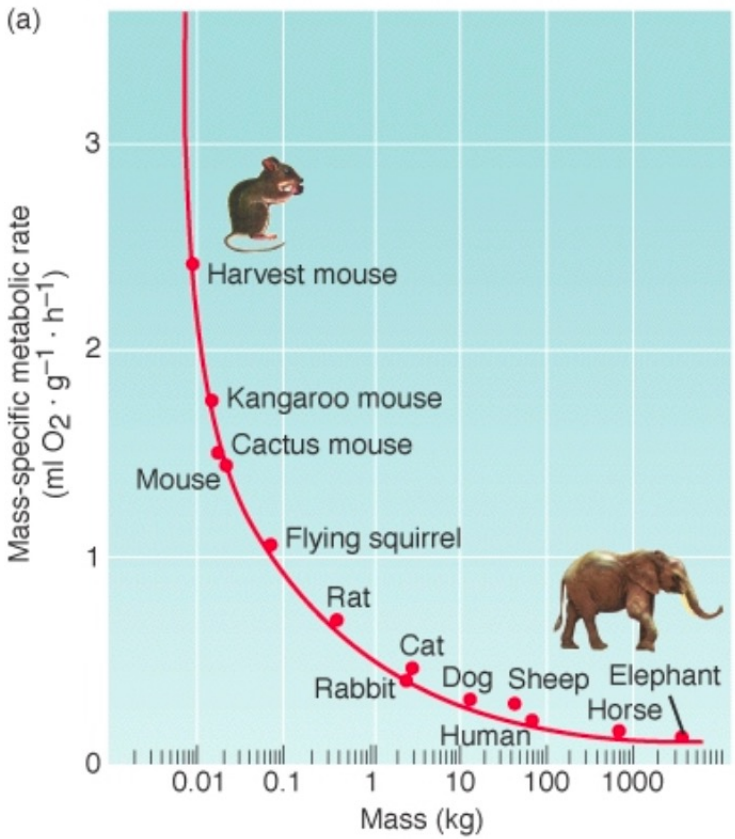
- a vector $v \in E$ is k-sparse if $[v]$ has exactly k non-zero values, and $k < \dim(E)$

On Moore Penrose Pseudo Inverse, Matrix Least Squares, Geometric Least Squares.



$$y_i = \beta_0 + \beta_1 x_i$$

$$y_i = \beta_0 + \beta_1 x_i + \epsilon_i$$



$$y_i = \beta_0 + \beta_1 x_i$$

$$y_i = \beta_0 + \beta_1 x_i + \epsilon_i$$

$$\begin{aligned} y_0 &= \beta_0 + \beta_1 x_0 + \epsilon_0 \\ y_1 &= \beta_0 + \beta_1 x_1 + \epsilon_1 \\ &\vdots \\ y_{N-1} &= \beta_0 + \beta_1 x_{N-1} + \epsilon_N - 1 \end{aligned}$$

$$Y = X\beta + \epsilon$$

$$e_i = (f(x_i) - y_i)^2$$

$$\begin{aligned} E &= \sum_{i=1}^{N} e_i \\ &= \sum_{i=1}^{N} (f(x_i) - y_i)^2 \\ &= \sum_{i=1}^{N} (\beta_0 + \beta_1 x_i - y_i)^2 . \end{aligned}$$

$$e_i = (f(x_i) - y_i)^2$$

$$\begin{aligned} E &= \sum_{i=1}^{N} e_i \\ &= \sum_{i=1}^{N} (f(x_i) - y_i)^2 \\ &= \sum_{i=1}^{N} (\beta_0 + \beta_1 x_i - y_i)^2 . \end{aligned}$$

$$\frac{\partial E}{\partial \beta_0} = 0 \quad \text{and} \quad \frac{\partial E}{\partial \beta_1} = 0,$$

$$e_i = (f(x_i) - y_i)^2$$

$$\begin{aligned} E &= \sum_{i=1}^{N} e_i \\ &= \sum_{i=1}^{N} (f(x_i) - y_i)^2 \\ &= \sum_{i=1}^{N} (\beta_0 + \beta_1 x_i - y_i)^2 . \end{aligned}$$

$$\frac{\partial E}{\partial \beta_0} = 0 \quad \text{and} \quad \frac{\partial E}{\partial \beta_1} = 0,$$

$$\begin{aligned}\frac{\partial E}{\partial \beta_1} &= \sum_{i=1}^{N} 2x_i \left(\beta_0 + \beta_1 x_i - y_i\right) \\ &= \sum_{i=1}^{N}(2\beta_0 x_i) + \sum_{i=1}^{N}(2\beta_1 x_i^2) - \sum_{i=1}^{N}(2x_i y_i) = 0\end{aligned}$$

and

$$\begin{aligned}\frac{\partial E}{\partial \beta_0} &= \sum_{i=1}^{N} 2 \left(\beta_0 + \beta_1 x_i - y_i\right) \\ &= \sum_{i=1}^{N}(2\beta_0) + \sum_{i=1}^{N}(2\beta_1 x_i) - \sum_{i=1}^{N}(2y_i) = 0.\end{aligned}$$

From the above two equations we have:

$$\begin{aligned}\sum_{i=1}^{N}(x_i y_i) &= \sum_{i=1}^{N}(\beta_0 x_i) + \sum_{i=1}^{N}(\beta_1 x_i^2) \\ \sum_{i=1}^{N}(y_i) &= \sum_{i=1}^{N}(\beta_0) + \sum_{i=1}^{N}(\beta_1 x_i)\end{aligned}$$

again, which can be written in matrix form as:

$$\begin{pmatrix} \sum_{i=1}^{N}(x_i y_i) \\ \sum_{i=1}^{N}(y_i) \end{pmatrix} = \begin{pmatrix} \sum_{i=1}^{N} x_i & \sum_{i=1}^{N} x_i^2 \\ \sum_{i=1}^{N} 1 & \sum_{i=1}^{N} x_i \end{pmatrix} \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix}$$

and then

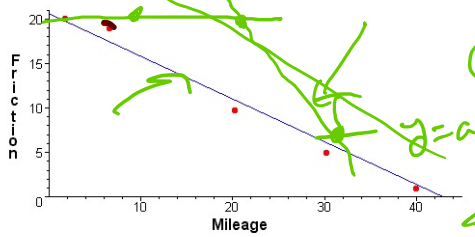
$$\begin{pmatrix} b_0 \\ b_1 \end{pmatrix} = \begin{pmatrix} \sum_{i=1}^{N} x_i & \sum_{i=1}^{N} x_i^2 \\ \sum_{i=1}^{N} 1 & \sum_{i=1}^{N} x_i \end{pmatrix}^{-1} \begin{pmatrix} \sum_{i=1}^{N}(x_i y_i) \\ \sum_{i=1}^{N}(y_i) \end{pmatrix}.$$

Example

An engineer is tracking the friction index over mileage of a breaking system of a vehicle. She expects that the mileage-friction relationship is approximately linear. She collects five data points that are show in the table below.

Mileage	2000	6000	20,000	30,000	40,000
Friction Index	20	18	10	6	2

The graph below shows these points



We are interested in the line that best fits the data. More specifically, if **b** is the vector of friction index data values and **y** is the vector consisting of y values when we plug in the mileage data for x and find y by the equation of the line, then we want the line that minimizes the distance between **b** and **y**. If the equation of the line is

$$ax + b = y$$

then we get the five equations

$$2a + b = 20$$
$$6a + b = 18$$
$$20a + b = 10$$
$$30a + b = 6$$
$$40a + b = 2$$

The corresponding matrix equation is

$$A\mathbf{x} = \mathbf{b}$$

or

$$\begin{pmatrix} 2 & 1 \\ 6 & 1 \\ 20 & 1 \\ 30 & 1 \\ 40 & 1 \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} 20 \\ 18 \\ 10 \\ 6 \\ 2 \end{pmatrix}$$

Although this does not have an exact solution, it does have a closest solution. We have

$$\begin{pmatrix} a \\ b \end{pmatrix} = (A^T A)^{-1} A^T \mathbf{b} = \begin{pmatrix} -0.48 \\ 20.6 \end{pmatrix}$$

We can conclude that the equation of the regression line is

$$y = -0.48x + 20.6$$

Least Squares

Definition and Derivations

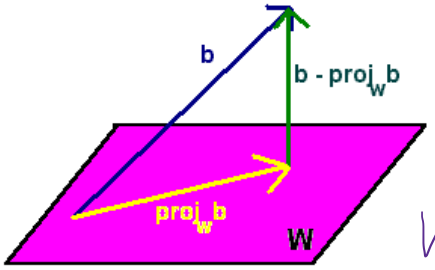
We have already spent much time finding solutions to

$$A\mathbf{x} = \mathbf{b}$$

If there isn't a solution, we attempt to seek the **x** that gets closest to being a solution. The closest such vector will be the **x** such that

$$A\mathbf{x} = \text{proj}_W\mathbf{b}$$

where W is the column space of A.



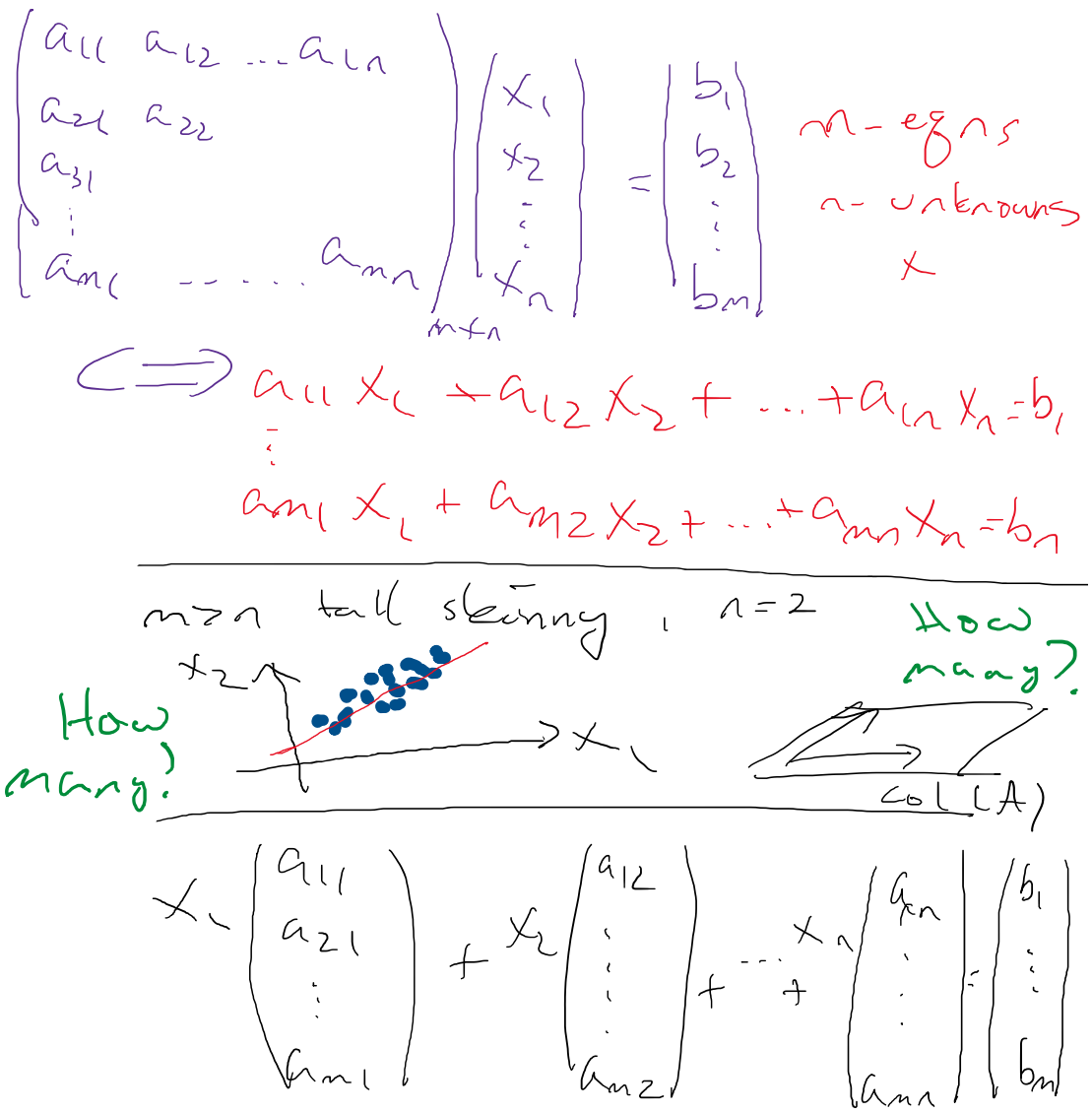
Notice that $\mathbf{b} - \text{proj}_W\mathbf{b}$ is in the orthogonal complement of W hence in the null space of A^T. Hence if **x** is a this closest vector, then

$$A^T(\mathbf{b} - A\mathbf{x}) = 0 \qquad A^TA\mathbf{x} = A^T\mathbf{b}$$

Now we need to show that A^TA nonsingular so that we can solve for **x**.

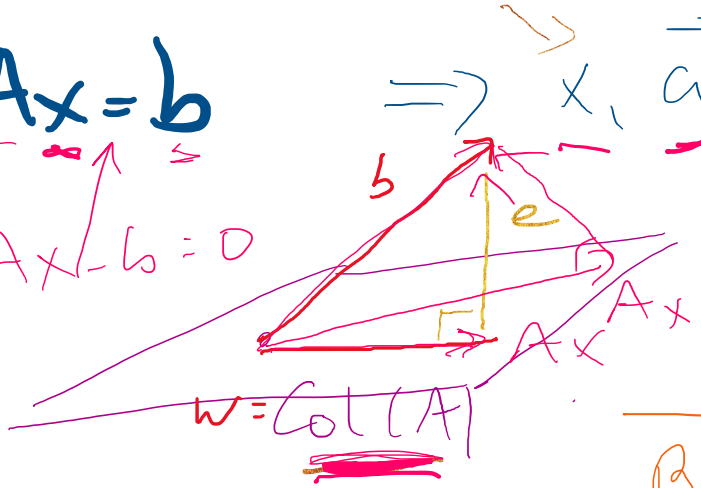
Lemma

If A is an m x n matrix of rank n, then A^TA is nonsingular.



$Ax = b$

$\Rightarrow x_1 \vec{a}_1 + x_2 \vec{a}_2 + \cdots + x_n \vec{a}_n = \vec{b}$



$A\tilde{x} = \text{proj}_W b$

$\tilde{e} = A\tilde{x} - b$

$Ax - b = 0$

$\tilde{x} = \underset{x}{\text{argmin}} \|Ax - b\|_2^2$

$\Rightarrow$ LS solution

Recall $u \perp v$ iff $u \cdot v = u^T v = 0$

$u \cdot v = \|u\| \|v\| \cos\theta$

$\Rightarrow A^T(A\tilde{x} - b) = 0 \iff$

$\begin{pmatrix} -a_1^T- \\ -a_2^T- \\ \vdots \\ -a_n^T- \end{pmatrix} \tilde{e} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$

$(Ax - b) = 0$ «

$A^T(Ax - b) = 0$

$\Rightarrow \vec{a}_i \perp (A\tilde{x} - b) \; \forall i$

$\Rightarrow (A\tilde{x} - b) \perp$ every vector in $\text{Col}(A) \iff$ Solve "normal eqns"

$A^T A \tilde{x}$

Theorem

Let A be an m x n matrix or rank n, then the system

$$A\mathbf{x} = \mathbf{b}$$

has the unique *least squares* solution

$$\mathbf{x} = (A^TA)^{-1}A^T\mathbf{b}$$

Example

Find the least squares solution to

$$A\mathbf{x} = \mathbf{b}$$

with

$$A = \begin{pmatrix} 1 & 3 \\ 2 & 4 \\ 1 & 6 \end{pmatrix} \quad \mathbf{b} = \begin{pmatrix} 4 \\ 1 \\ 3 \end{pmatrix}$$

Solution

We can quickly check that A has rank 2 (the first two rows are not multiples of each other). Hence we can compute

$$\mathbf{x} = (A^TA)^{-1}A^T\mathbf{b} = \begin{pmatrix} -0.377 \\ .662 \end{pmatrix}$$

Notice that

$$A\mathbf{x} = \begin{pmatrix} 1.61 \\ 1.90 \\ 3.60 \end{pmatrix}$$

not exactly **b**, but as close as we are going to get.

A^TA - never form that

$Ax = b$ $m \times n$

- Square & consing. A^{-1} d.ne. $x = A^{-1}b$ unique b.
- Not, $m > n$

$A^TAx = A^Tb$ normal eqns

Covariance matrix ; (if identical).

$m > n$ - tall skinny

Best Fitting Curves

Often, a line is not the best model for the data. Fortunately the same technique works if we want to use other nonlinear curves to fit the data. Here we will explain how to find the least squares cubic. The process for other polynomials is similar.

Example

A bioengineer is studying the growth of a genetically engineered bacteria culture and suspects that is it approximately follows a cubic model. He collects six data points listed below

Time in Days	1	2	3	4	5	6
Grams	2.1	3.5	4.2	3.1	4.4	6.8

He assumes the equation has the form

$$ax^3 + bx^2 + cx + d = y$$

This gives six equations with four unknowns

$$\begin{aligned} a + b + c + d &= 2.1 \\ 8a + 4b + 2c + d &= 3.5 \\ 27a + 9b + 3c + d &= 4.2 \\ 64a + 16b + 4c + d &= 3.1 \\ 125a + 25b + 5c + d &= 4.4 \\ 216a + 36b + 6c + d &= 6.8 \end{aligned}$$

The corresponding matrix equation is

$$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 8 & 4 & 2 & 1 \\ 27 & 9 & 3 & 1 \\ 64 & 16 & 4 & 1 \\ 125 & 25 & 5 & 1 \\ 216 & 36 & 6 & 1 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} 2.1 \\ 3.5 \\ 4.2 \\ 3.1 \\ 4.4 \\ 6.8 \end{pmatrix}$$

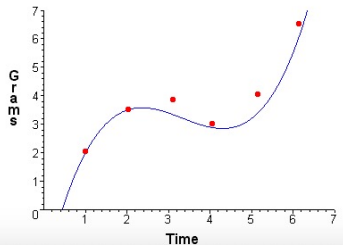
We can use the least squares equation to find the best solution

$$\begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = (A^T A)^{-1} A^T \mathbf{b} = \begin{pmatrix} 0.2 \\ -2.0 \\ 6.1 \\ -2.3 \end{pmatrix}$$

So that the best fitting cubic is

$$y = 0.2x^3 - 2.0x^2 + 6.1x - 2.3$$

The graph is shown below



Matrix Formulation of LS slick for general Models

Grams = f(time)
y = f(x)

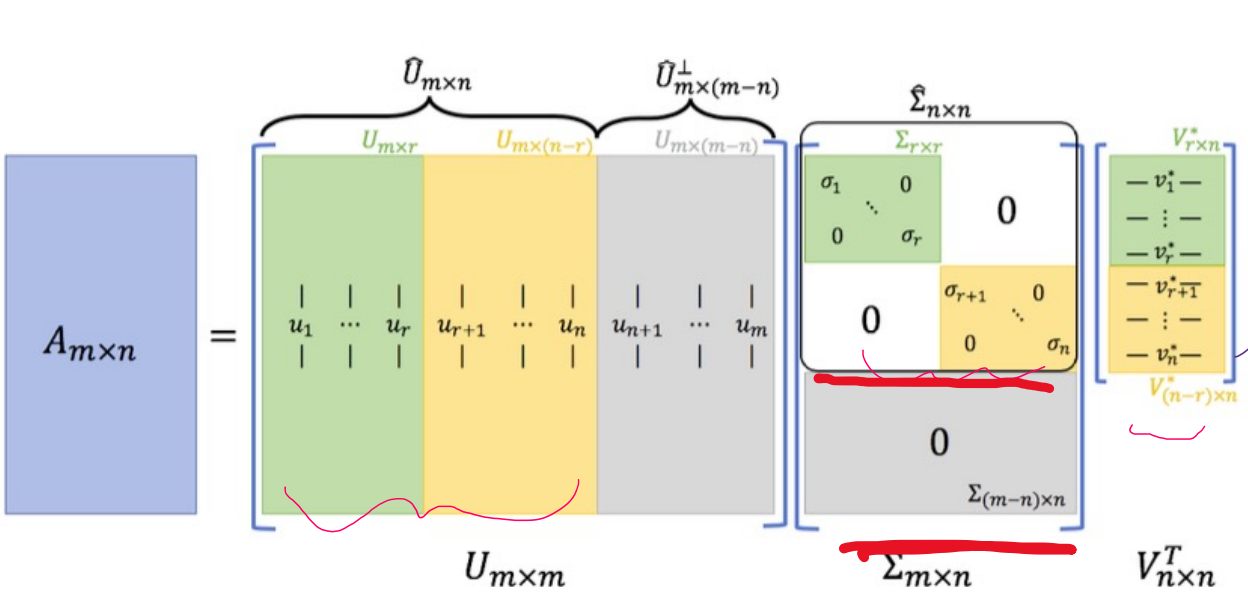
$$a \cdot 1 = a \cdot 1^3 + b \cdot 1^2 + c \cdot 1 + d \qquad i = 1$$
$$3.5 = a \cdot 2^3 + b \cdot 2^2 + c \cdot 2 + d \qquad i = 2$$

$$Ax = b$$
$$X\beta = y$$

$$\begin{pmatrix} x_1^3 & x_1^2 & x_1 & 1 \\ x_2^3 & x_2^2 & x_2 & 1 \\ \vdots & \vdots & \vdots & \vdots \\ x_n^3 & x_n^2 & x_n & 1 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix}$$

X β $=$ y

A x $=$ b



$$A^T A x = A^T b$$

$$A_{m\times n} = \hat{U}_{m\times n}\hat{\Sigma}_{n\times n}\hat{V}^T_{n\times n}$$

$$= \begin{bmatrix} | & | & | & | \\ u_1 & u_2 & \dots & u_n \\ | & | & | & | \end{bmatrix} \begin{bmatrix} \sigma_1 & 0 & \dots & 0 \\ 0 & \sigma_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sigma_n \end{bmatrix} \begin{bmatrix} - & v_1^T & - \\ - & v_2^T & - \\ - & \vdots & - \\ - & v_n^T & - \end{bmatrix}$$

$$\hat{V}\hat{\Sigma}^T\hat{U}^T\tilde{b}$$

$$\hat{V}\hat{\Sigma}^T\hat{\Sigma}\hat{V}^T\tilde{x} = \hat{V}\hat{\Sigma}^T\hat{U}^T\tilde{b}$$

$$\left(\hat{\Sigma}^T\hat{\Sigma}\right)\hat{V}^T\tilde{x} = \hat{\Sigma}^T\hat{U}^T\tilde{b}$$

$$\hat{V}^T\tilde{x} = \left(\hat{\Sigma}^T\hat{\Sigma}\right)^{-1}\hat{\Sigma}^T\hat{U}^T\tilde{b}$$

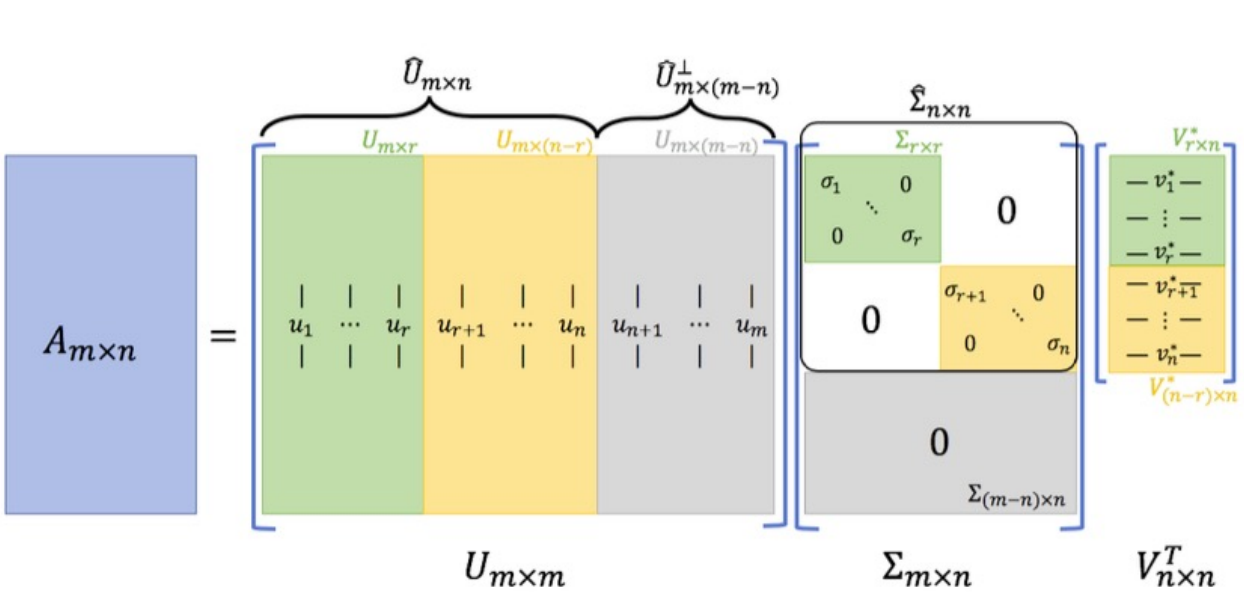
$$\tilde{x} = \hat{V}\left(\hat{\Sigma}^T\hat{\Sigma}\right)^{-1}\hat{\Sigma}^T\hat{U}^T\tilde{b}$$

A^+

$\hat{\Sigma}$

$\hat{\Sigma}^T\hat{\Sigma}$ ~ when exist

$\sigma_0 = 0$



$$\hat{\Sigma}^{+} = (\hat{\Sigma}^T\hat{\Sigma})^{-1}\hat{\Sigma}$$

if $\hat{\Sigma}$ is invertible

- if $r = n$
  if $\sigma_i > 0$, $i < n$

But if $r < n$

- $\sigma_r > 0$, $\sigma_{r+1} = 0$

$$(\hat{\Sigma}^T\hat{\Sigma}) = \begin{pmatrix} 1/\sigma_1^2 & & & 0 \\ & \ddots & & \\ & & 1/\sigma_r^2 & \\ 0 & & & 0 \end{pmatrix}$$

$\frac{1}{\sigma}$:= 0

$$A_{m\times n} = \hat{U}_{m\times n}\hat{\Sigma}_{n\times n}\hat{V}^T_{n\times n}$$

$$= \begin{bmatrix} | & | & | & | \\ u_1 & u_2 & \dots & u_n \\ | & | & | & | \end{bmatrix} \begin{bmatrix} \sigma_1 & 0 & \dots & 0 \\ 0 & \sigma_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sigma_n \end{bmatrix} \begin{bmatrix} - & v_1^T & - \\ - & v_2^T & - \\ - & \vdots & - \\ - & v_n^T & - \end{bmatrix}$$

LS soln = solve normal equations

$$A^T A \tilde{x} = A^T b$$

When inverse exists

$$\tilde{x} = (A^T A)^{-1} A^T b$$

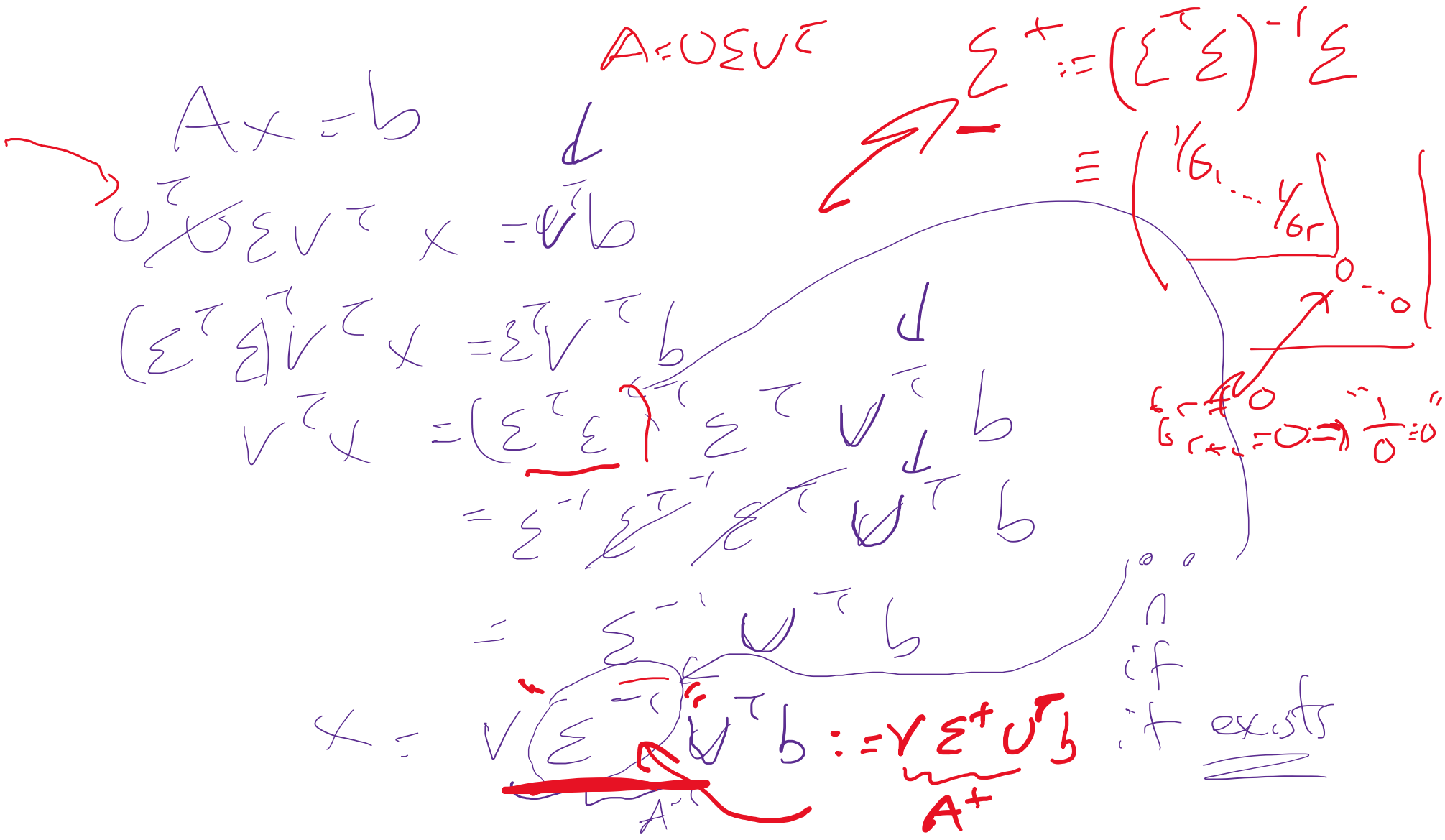
$$\equiv A^+ b$$

Moore - Penrose Pseudo-Inverse

A^+

$Ax = b$

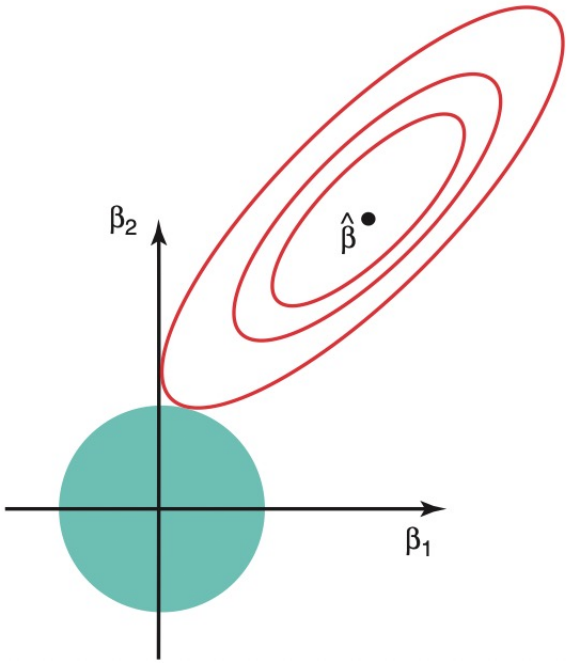
- In terms of SVD?
- and what if inverse doesn't exist.


$A = U \Sigma V^T$
$\Sigma^+ := (\Sigma^T \Sigma)^{-1} \Sigma$
$\equiv \begin{pmatrix} 1/\sigma_1 & & & \\ & \ddots & & \\ & & 1/\sigma_r & \\ & & & 0 \cdots 0 \end{pmatrix}$
$Ax = b$
$U^T U \Sigma V^T x = U^T b$
$(\Sigma^T \Sigma) V^T x = \Sigma^T U^T b$
$V^T x = (\Sigma^T \Sigma)^{-1} \Sigma^T U^T b$
$= \Sigma^{-1} \Sigma^{-T} \Sigma^T U^T b$
$= \Sigma^{-1} U^T b$
$x = V \Sigma^{-1} U^T b := V \Sigma^+ U^T b$
A^+
if
if exists

Tikhonov Regularization for inverse problems and preventing over fitting

-Ridge regression

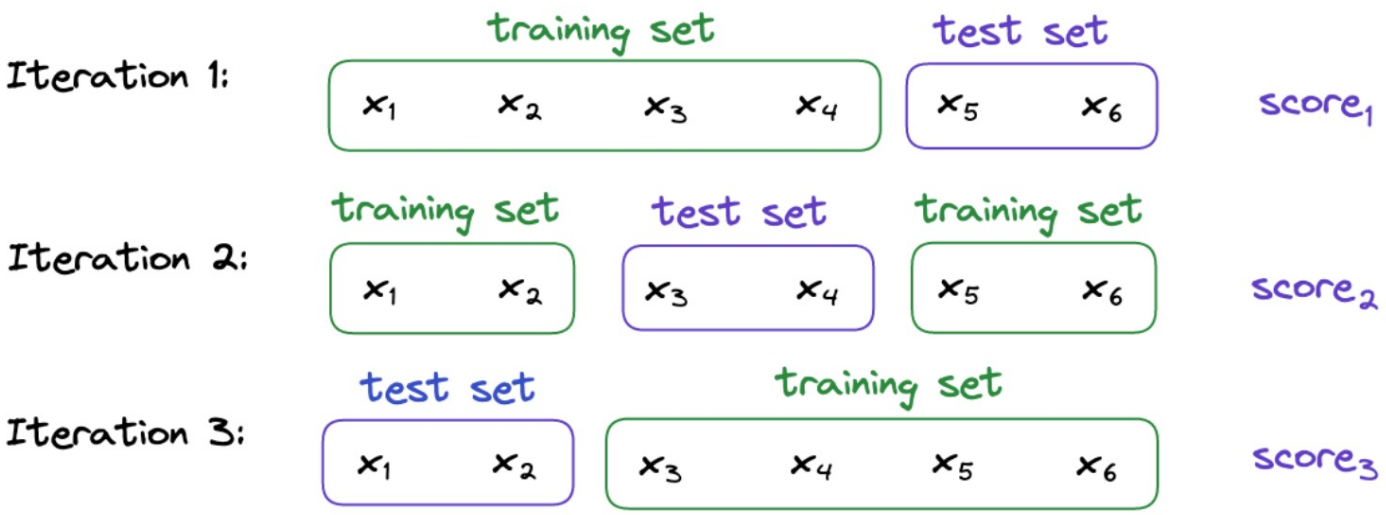
-Lasso regression



K-fold leave 1 out cross validation – A version cross validation

Then average scores of testing out of training sample

K-fold leave l out cross validation



Compressed sensing and L1 regularization – Sparse Signal Recovery

Exact sparsity vs. approximate

$$f(x) = \sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \frac{x^9}{9!} - \frac{x^{11}}{11!} + \frac{x^{13}}{13!} + \cdots$$

$f(x)$ — $f_9(x)$ — $N = 9$

x

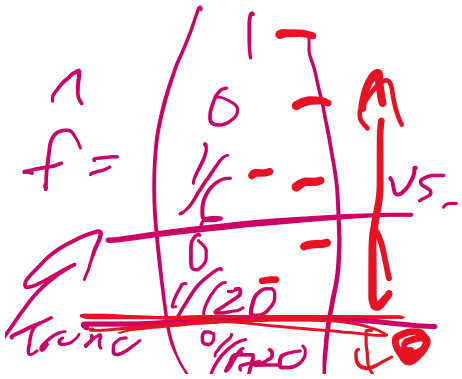
$\phi_n(x) = x^n$

$$\sin x = \sum_{n=1}^{\infty} a_n \phi_n(x) \simeq \sum_{n=1}^{N} a_n \phi_n(x)$$

$f(x)$ = $f_N(x)$

$\| f - f_N(x) \|$

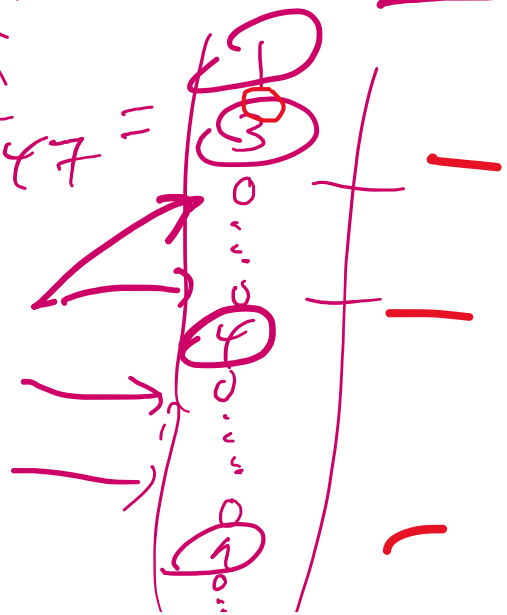
$= \| (\sin x - (x - \frac{x^3}{3!}) \quad \cdots \quad) \|$



$\hat{f} = \begin{pmatrix} 0 \\ 1/6 \\ 0 \\ \vdots \end{pmatrix}$ $N = 3$

- Truncation imposes sparsity.
  (just truncate after a large number.)

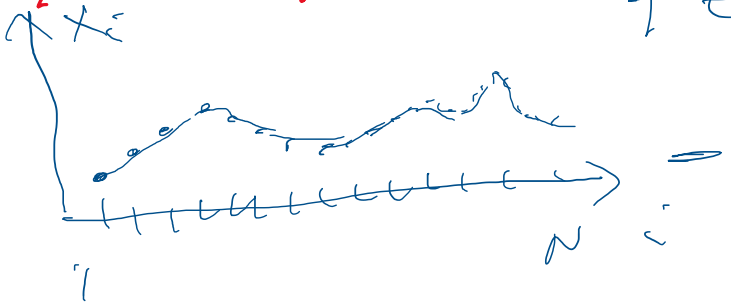
- $f(x) = \sin x + 3\sin 3x + 9\sin 12x + 9\sin 47x$

$\hat{f} = \hat{f}_{47} =$ 

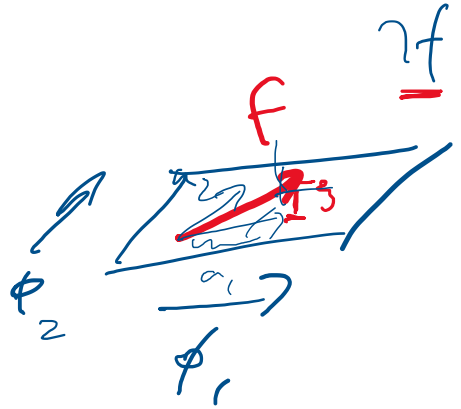
just 1, 3, 4, 9 and tell you where to put them.

3 Truncation as Projection
by matrix Mult.

$f \in \mathcal{H}$



$=$



(say from $f(\bar{s}) = \bar{s}^3$)

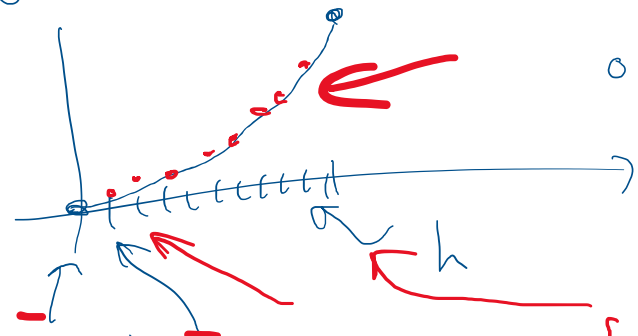
$x_i := f(s_i)$

N

$x_i = [x]_i$

$[x]_3 = x_3 = 0$

$0 \le \bar{s} \le 2$

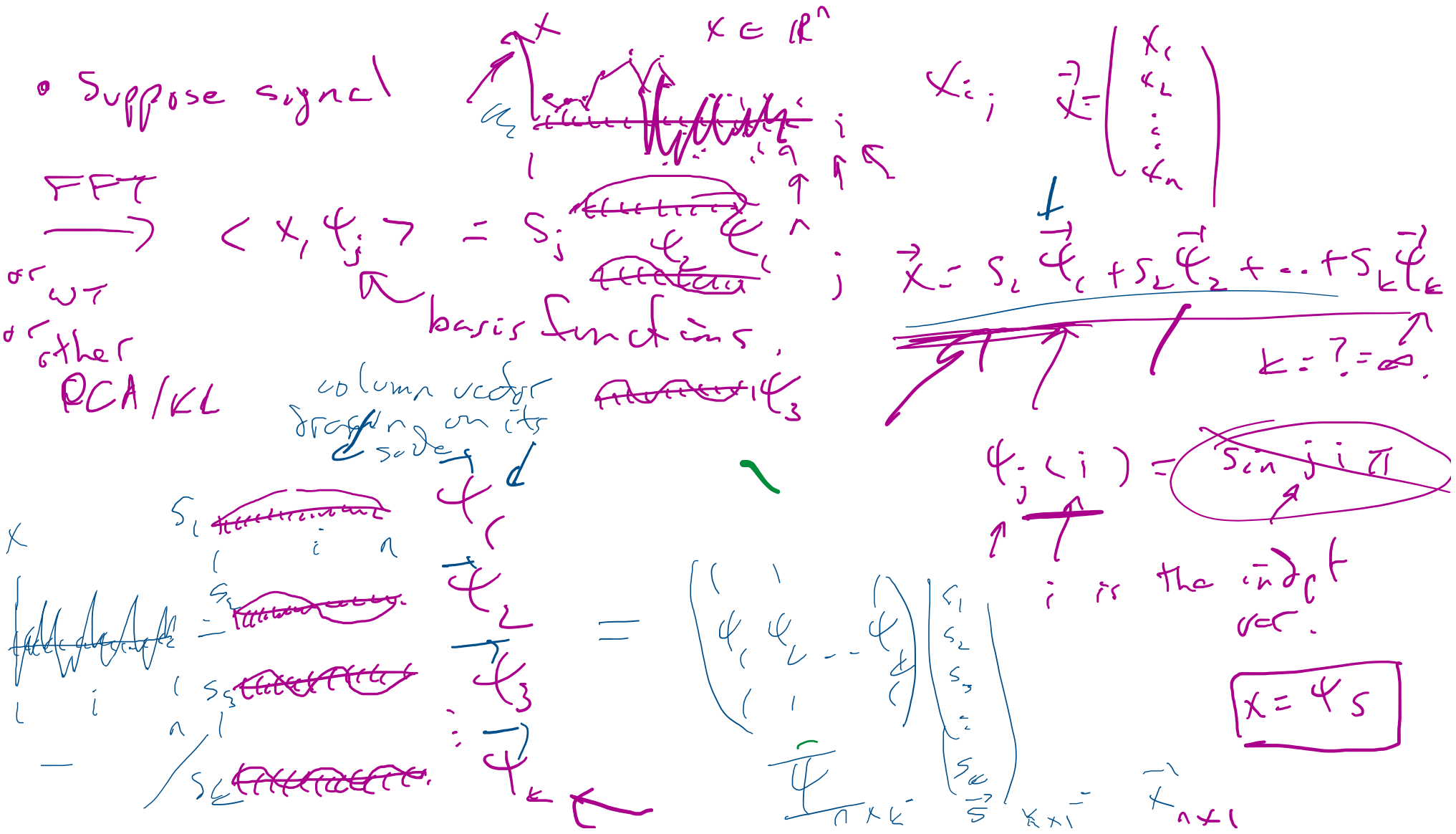


$N = 100$, Uniform

$h := \frac{2-0}{N} = \frac{2}{100} = \frac{1}{50}$

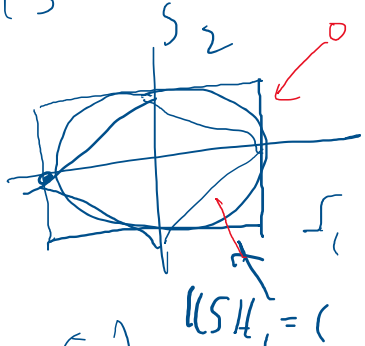
$\vec{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ \vdots \end{pmatrix} = \begin{pmatrix} 0 \\ (\frac{1}{50})^3 \\ (\frac{2}{50})^3 \\ \vdots \end{pmatrix}$

$s_1 = 0$ $s_2 = \frac{1}{50}$ $s_3 = \frac{2}{50}$ $s_{100} = 2$


Suppose signal
$x \in \mathbb{R}^n$
FFT
or WT
or other
PCA/KL
$\langle x, \psi_j \rangle = s_j$
basis functions
column vectors
$\vec{x} = s_1 \vec{\psi}_1 + s_2 \vec{\psi}_2 + \dots + s_k \vec{\psi}_k$
$k = ? = \infty$
$\psi_j(i) = \sin j i \pi$
i is the index
$\vec{x} = \Psi \vec{s}$
$n \times k$
$k \times 1$
$n \times 1$

$\|s\|_0 = \#$ of not zeros.

"Balls" $\|s\|_p = 1$



$\|s\|_1 = |s_1| + |s_2| + \dots + |s_n|$

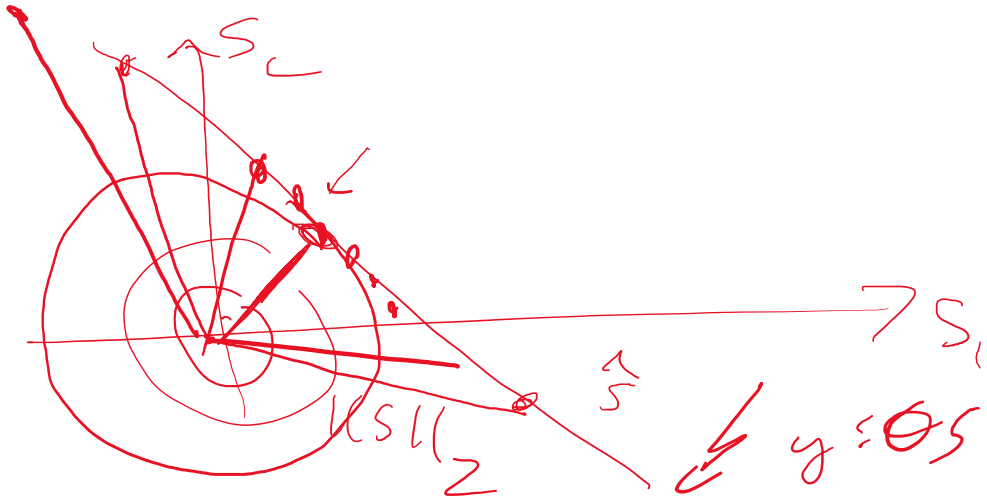
$\|s\|_2 = (s_1^2 + s_2^2 + \dots + s_n^2)^{1/2}$ Euclidean

$\|s\|_\infty = \rightarrow \max |s_i|$

$\|\vec{s}\|_p = (s_1^p + s_2^p + \dots + s_n^p)^{1/p}$

$s \in \mathbb{R}^n$

Want to solve $\|w\|_0 \leq 10$ — But have to test all subsets NP-complete.



$y = \Theta s$

$\rho = 1$

$$s = \underset{\text{subj}}{\operatorname{argmin}} \|\hat{s}\|_2 \quad \text{on the circle.}$$

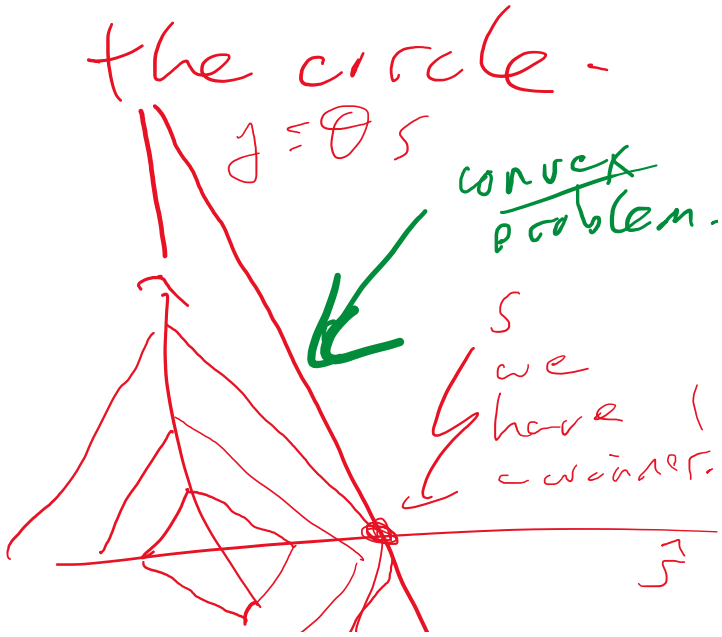
$y = \Theta \hat{s}$

Basis pursuit — Convex optimization.

convex problem.

$$s = \underset{\text{subj}}{\operatorname{argmin}} \|\hat{s}\|_1$$

$\Rightarrow y = \Theta \hat{s}$



$$\underset{n\times 1}{\vec{x}} = \underset{n\times n}{\Psi}\,\underset{n\times 1}{\vec{s}} = \begin{pmatrix} | & | & & | \\ \psi_1 & \psi_2 & \cdots & \psi_n \\ | & | & & | \end{pmatrix}_{n\times n} \begin{pmatrix} s_1 \\ s_2 \\ \vdots \\ s_n \end{pmatrix}_{n\times 1} \qquad 3.1$$

$\vec{x} \in \text{col}(\Psi)$

= linear combo of columns

$n \times 1$ column

$$= s_1\psi_1 + s_2\psi_2 + \cdots + s_n\psi_n$$

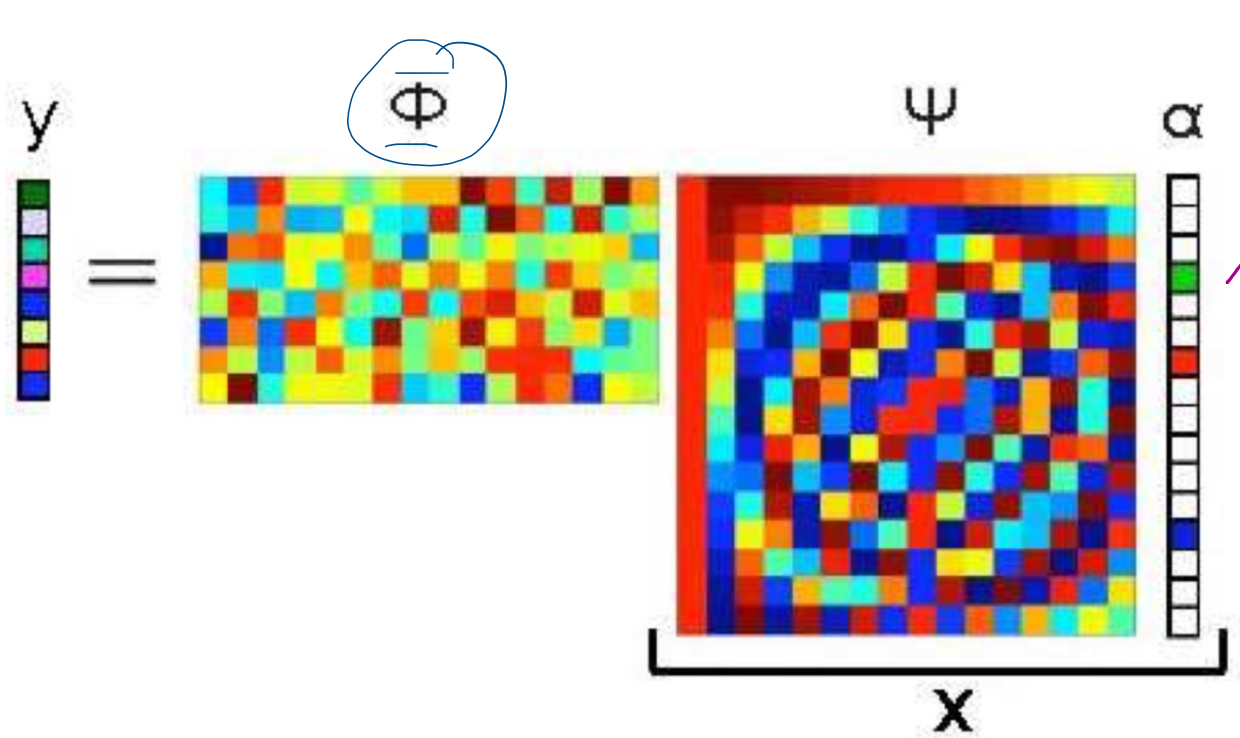
(linear combo of columns)

$$\vec{x} = \sum_{i=1}^{N} s_i\psi_i \approx \sum_{i=1}^{n} s_i\psi_i \qquad n < N$$

$$\vec{x} = \left(\psi_1\ \psi_2\ \cdots\ \psi_n \mid \psi_{n+1}\ \cdots\ \psi_N\right)\begin{pmatrix} s_1 \\ s_2 \\ \vdots \\ s_n \\ \hline s_{n+1} \\ \vdots \\ s_N \end{pmatrix} \qquad s_{n+1} = \cdots = s_N = 0$$

$$\Rightarrow \Psi_N \text{ vs. } \Psi_n \qquad \Psi_n = (\quad \Psi_N)$$

Random measurements can be used for signals sparse in any basis.



$x = \Psi S$; $\Psi_{n \times m}$

$S_{m \times 1}$

$x_{n \times 1}$

- choose $k = m$ modes
- a sampler of space for $x \& \Psi$:

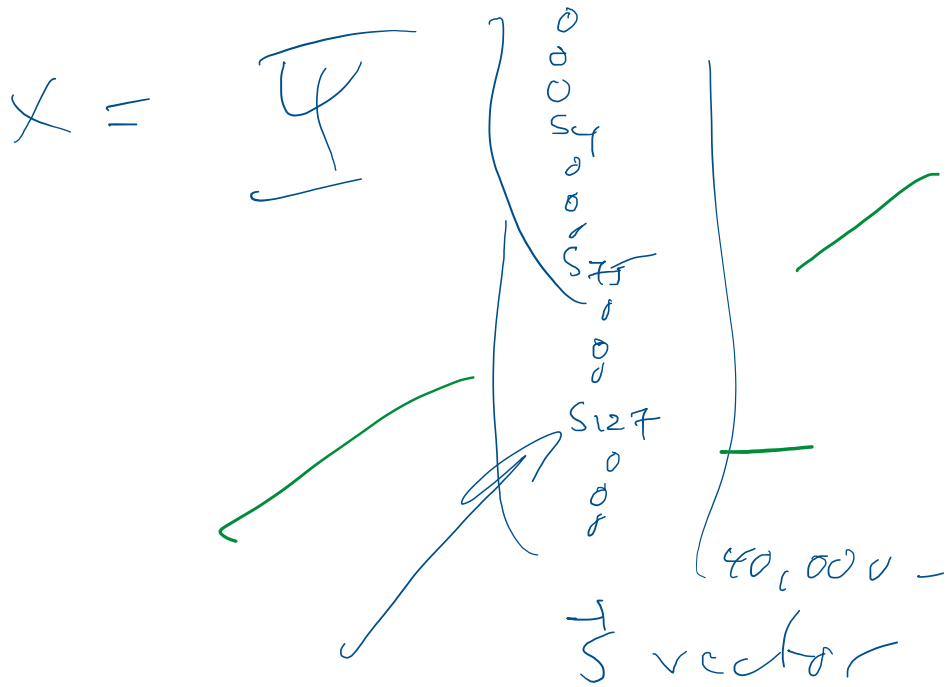
Shannon-Nyquist

- Choose $\downarrow$ at least twice as big as frequency content of x

- $\left| \hat{x} \right| = \left(\left| \Psi \right| \right)$ $\begin{matrix} s_1 \\ \vdots \\ s_k \\ \vdots \\ s_m \end{matrix}$ $\left.\begin{matrix} \\ \\ \\ \\ \end{matrix}\right\}$ 20,000 non zero

$\left.\begin{matrix} \\ \\ \\ \\ \end{matrix}\right\}$ 40,000

- if S is full (non zeros.) then I need at least to truncate "left".



s is
- k-sparse if s has exactly or more than k non zero values
- Call sparse if $k \ll m$

$3 \ll 40{,}000$

s_i is a scalar

$$x = \frac{1}{n}\sin 4it \, \psi_4 + \frac{1}{n}\sin 75it \, \psi_{75} + s_{127}\psi_{127}$$

$$= s_4\psi_4 + s_{75}\psi_{75} + s_{127}\psi_{127}$$

$$= \hat{s}_1\hat{\psi}_1 + \hat{s}_2\hat{\psi}_2 + \dots \hat{s}_m\hat{\psi}_m$$

3

- sparse basis but widely full in another basis.

- what if x is not sparse w.r.t. $B = \{\psi_i\}_{i=1}^{n}$
  it may be almost sparse w.r.t. B.

$$x = s_1\psi_1 + s_2\psi_2 + s_3\psi_3 + \cdots + s_n\psi_n$$

- and all or most $s_i \neq 0$ but

$$s_4 > s_{75} > s_{127} \gg s_i \approx \varepsilon$$

most of the $s_i \approx 0$.

$$s \approx \begin{pmatrix} 1.1 \cdot 10^{-10} \\ 1.10^{-8} \\ 3.10^{-7} \\ 14 \\ 2.10^{-6} \\ \vdots \\ 3.10^{-11} \\ 3 \end{pmatrix}$$

- $x = \Psi s$ ← s coefficients
  (signal) ← assume s is k-sparse or at least almost k-sparse.
- $\|s - \hat{s}\|_2 < \delta$ for $\delta > 0$ small & $\hat{s}$ is sparse

---

Observe $y = Cx$ ← projection matrix C.

y: $p \times 1$, C: $p \times n$, x: $n \times 1$ ← signal I want

$p < n$

- → $y_{p\times 1} = \Theta_{p\times n}\, s_{n\times 1}$ ✓

- ? $s = \Theta^{-1} y$ ~~(crossed out)~~ $p << n$

Solve for $s \in$ w/ s, $x = \Psi s$

singular problem ⇒ no unique solution

- inverse problems field say this ill-posed.

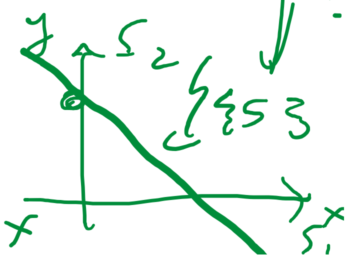
- Hadamard define a well posed problem as one where
  - solution exists
  - uniquely
  - and cts(ly) w.r.t. to the data.

else call ill-posed

- Extreme: what if $p = 1$ $(n = 2)$

Ex: $y_{1\times 1} = \begin{pmatrix} 1 & 1 \end{pmatrix}_{1\times 2} \begin{pmatrix} s_1 \\ s_2 \end{pmatrix}$ ⇒ $y = s_1 + s_2$

$s_2 = y - s_1$ / $y = 7 - x$



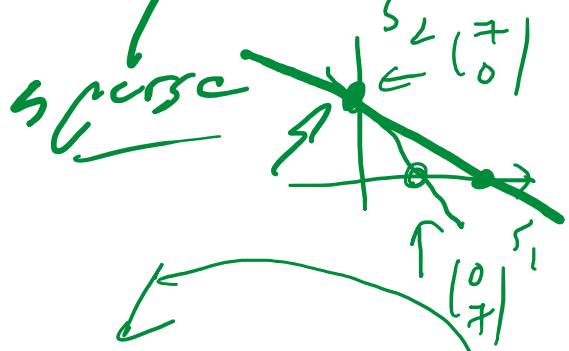
But $7 = \begin{pmatrix} 1 & 1 \end{pmatrix} \begin{pmatrix} s_1 \\ s_2 \end{pmatrix}$ (2)

$y = \Theta s$

$\rho = 1, n = 2$

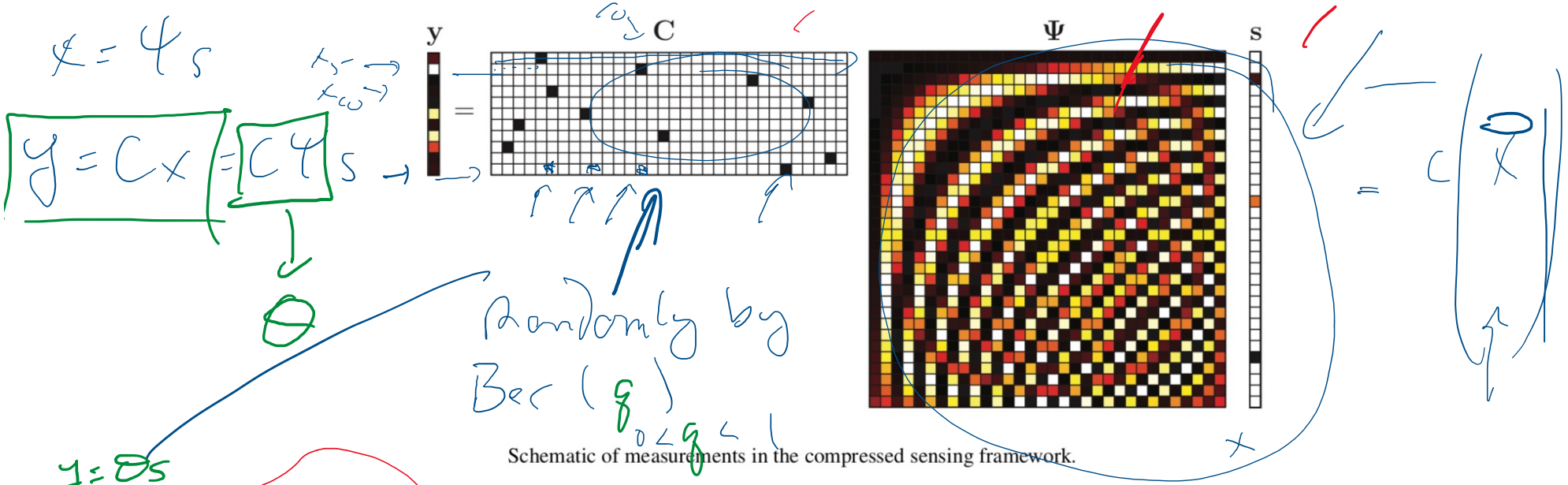
but true value were $s = \begin{pmatrix} 7 \\ 0 \end{pmatrix}$

if tell you tell me. and s is sparse.

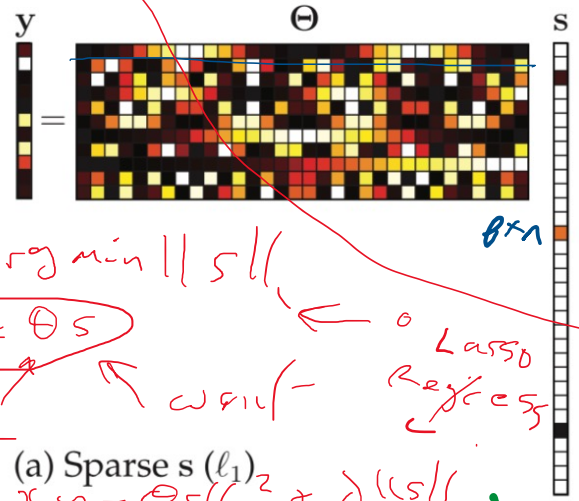


- let's try. $\vec{s} = \underset{\vec{s}}{\operatorname{argmin}} \|\vec{s}\|_0$ subject to $y = \Theta \vec{s}$

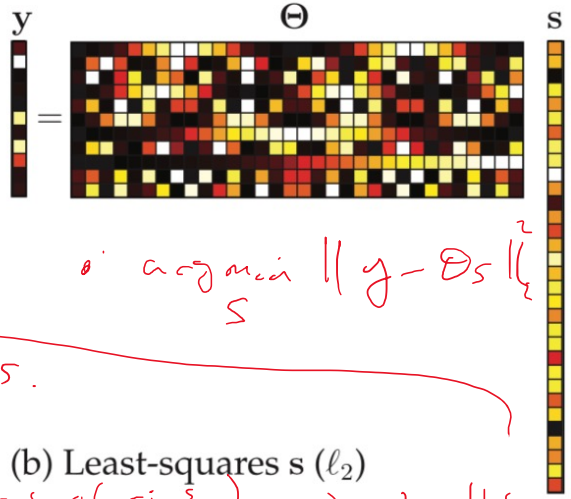
$\left\{ \begin{pmatrix} 7 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 7 \end{pmatrix} \right\}$



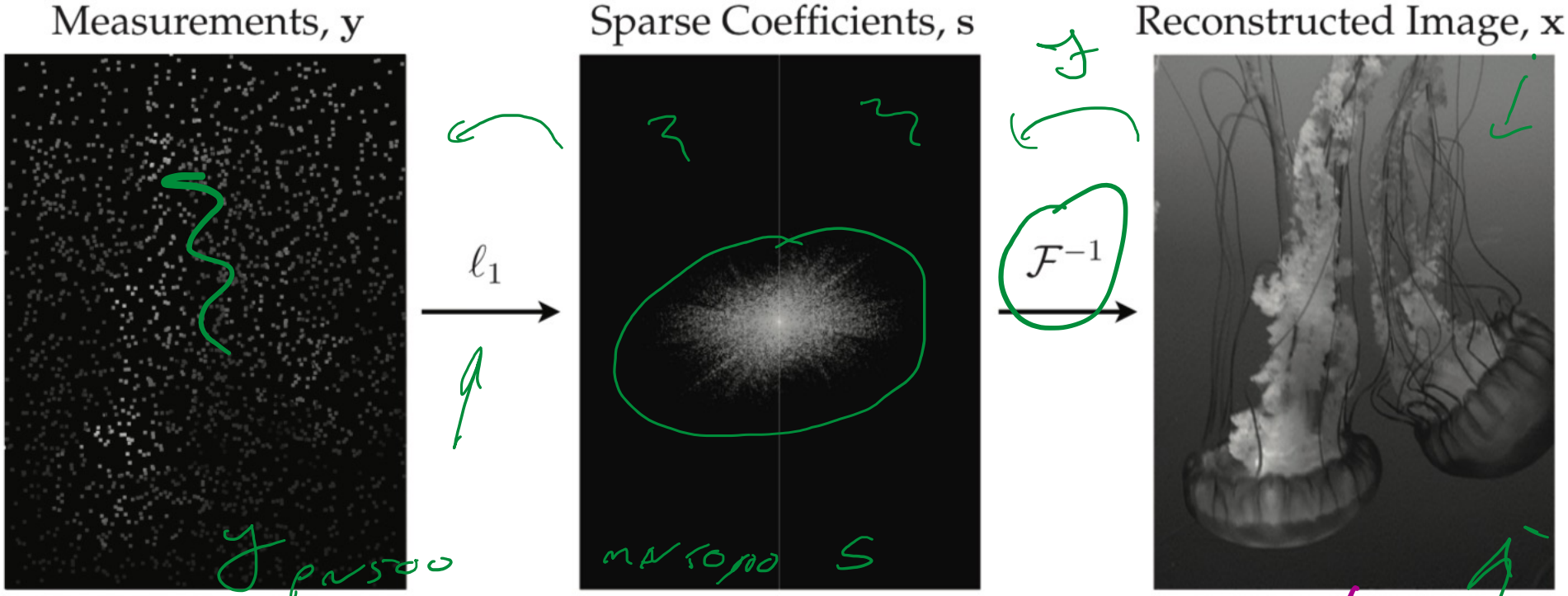
Schematic of measurements in the compressed sensing framework.



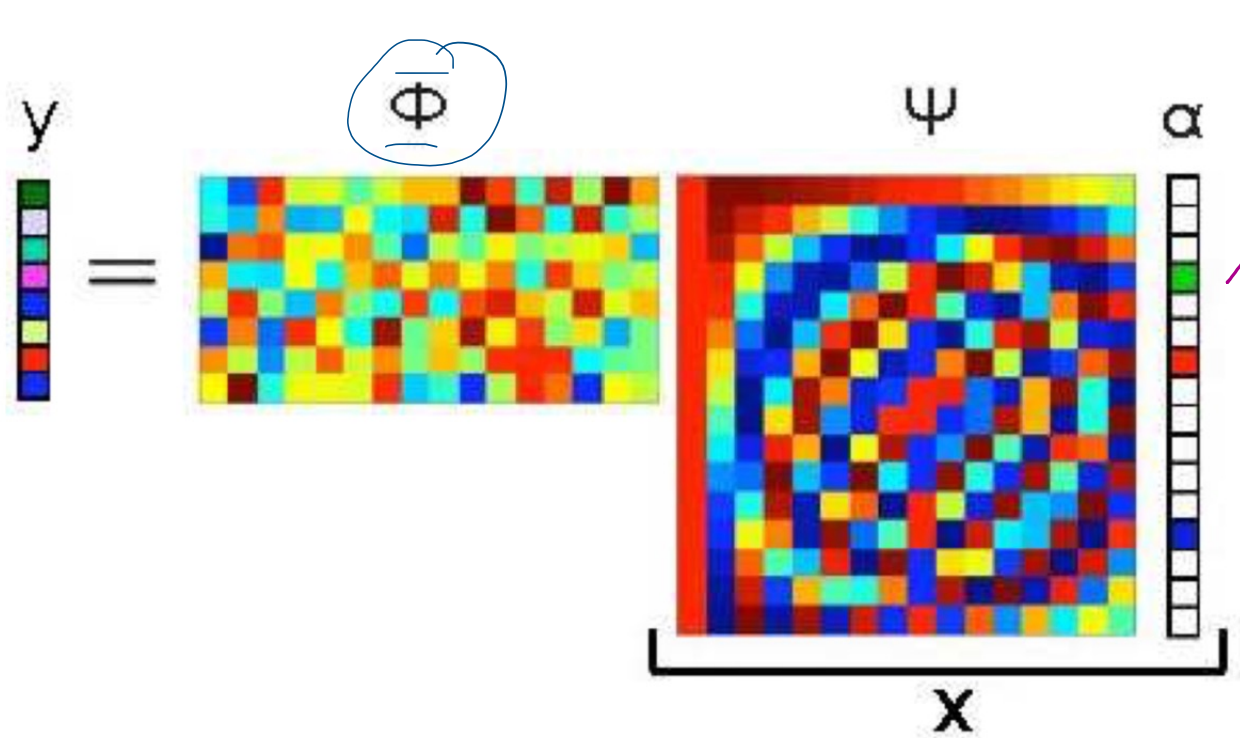
(a) Sparse s (ℓ_1)



(b) Least-squares s (ℓ_2)


Measurements, y
Sparse Coefficients, s
Reconstructed Image, x
ℓ_1
$\mathcal{F}^{-1}$

Random measurements can be used for signals sparse in any basis.

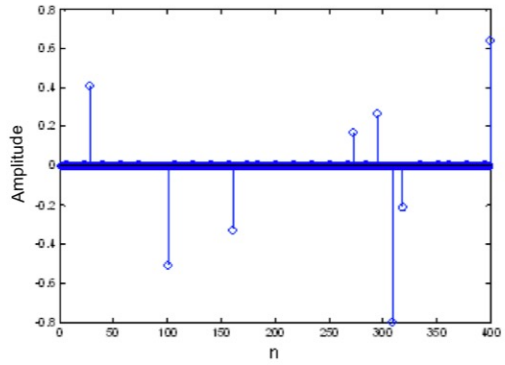


$$y_k =< \phi_k, x >; \quad k = 1, ..., M; \quad \text{with} \quad M \ll N$$

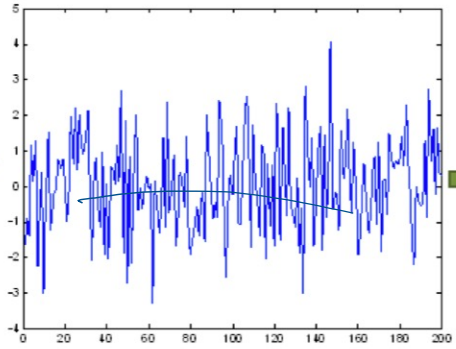
- Need to solve an under determined system of equations $y = \Phi x$.
- Infinitely solutions for the system since $M \ll N$.
- A sparse solution x is recovered from y by solving the following inverse problem

$$(P0): \quad \min_x \|x\|_{\ell_0} \quad s.t. \quad y = \Phi x$$

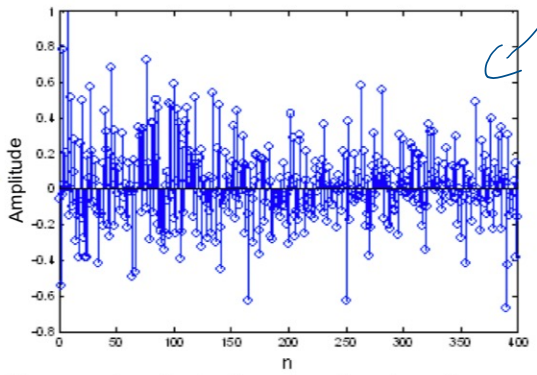
Example of the recovery of an under determined system of equations:



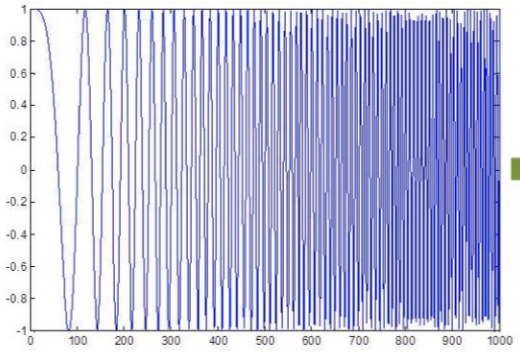
Original sparse signal



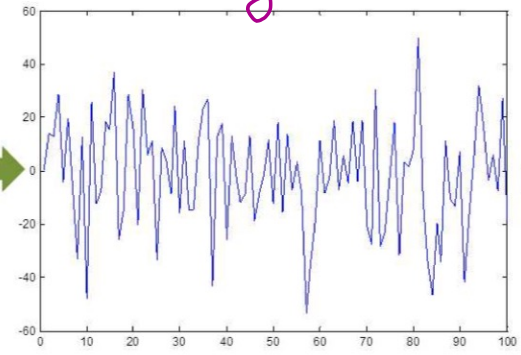
Compressed measurements



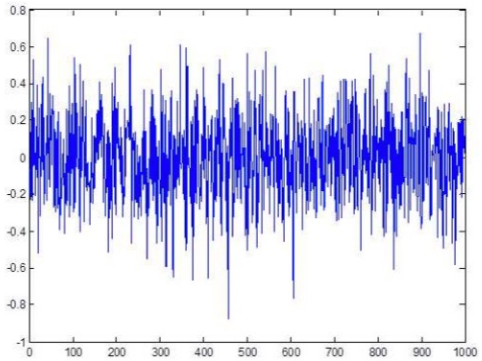
Reconstructed signal using least-squares.
Solution not sparse



Original sparse signal
in Time-Frequency basis



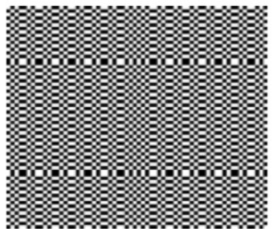
Compressed Measurements



Reconstructed signal using
least-squares

- Sparsity is what makes it possible to recover a signal from undersampled data.
- The number of measurements we need for successful reconstruction depends on the nature of the waveforms ϕ_k, and S

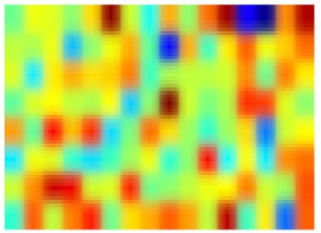
1. Incoherent Orthobasis



Ordered Hadamard Ensemble

Scrambled Block Hadamard Ensemble

2. Random waveforms ϕ_k



Gaussian Random Ensemble

Incoherent Orthobasis Example

Example of incoherent basis: the ”spike” basis (identity) and the Fourier basis.

Consider the case where the dictionary is the union of two orthobasis:

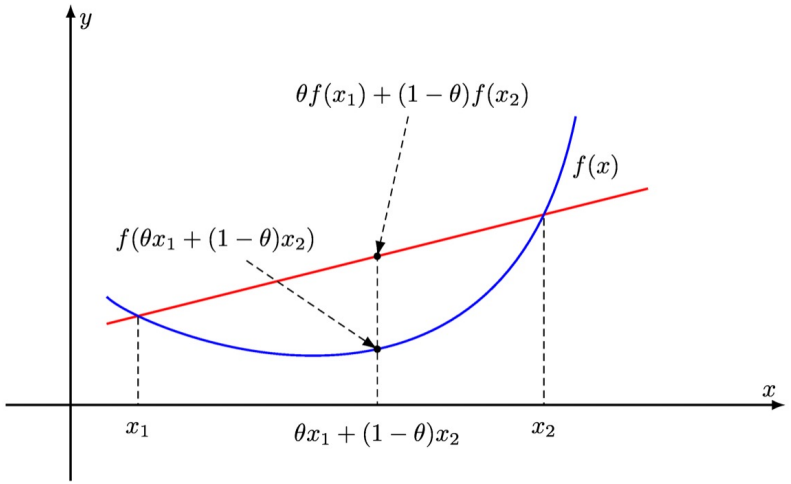
- I: the ”spike” basis (identity).
- F: the Fourier basis (sinusoids).

$$\Phi = [I; F]$$

where I is a $N \times N$ matrix and F is a $N \times N$ matrix with

$$f_{m,\ell} = \frac{1}{\sqrt{(N)}} e^{j2\pi(m-1)(\ell-1)/N}$$

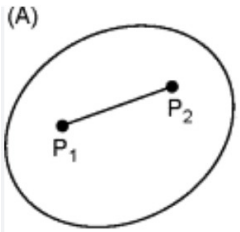
Convex Optimization – optimize a convex objective function over a convex domain

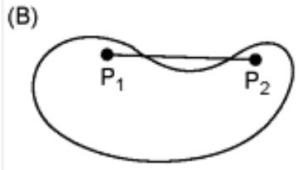


Convex function

$$f\left(tx_1 + (1-t)x_2\right) \leq tf\left(x_1\right) + (1-t)f\left(x_2\right)$$

Convex Domain





Compressive Sensing

[Candes, Romberg, Tao; Donoho]

- Signal x is K-*sparse* in basis/dictionary Ψ
  - WLOG assume sparse in space domain $\Psi = I$
- Replace samples with *few linear projections* $y = \Phi x$

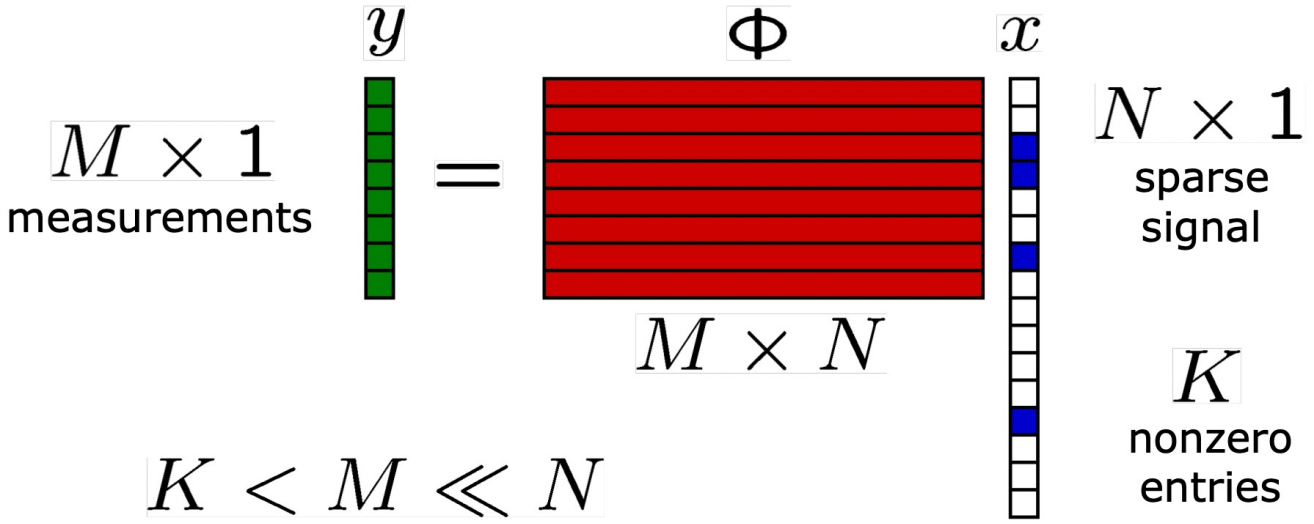
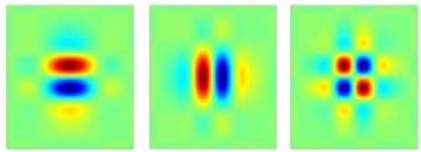


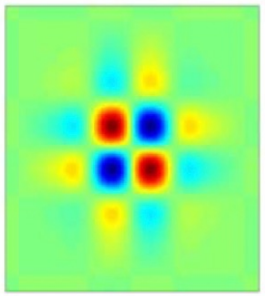
Image Representation

N
pixels

$$x = \sum_{i=1}^{N} \alpha_i \psi_i$$



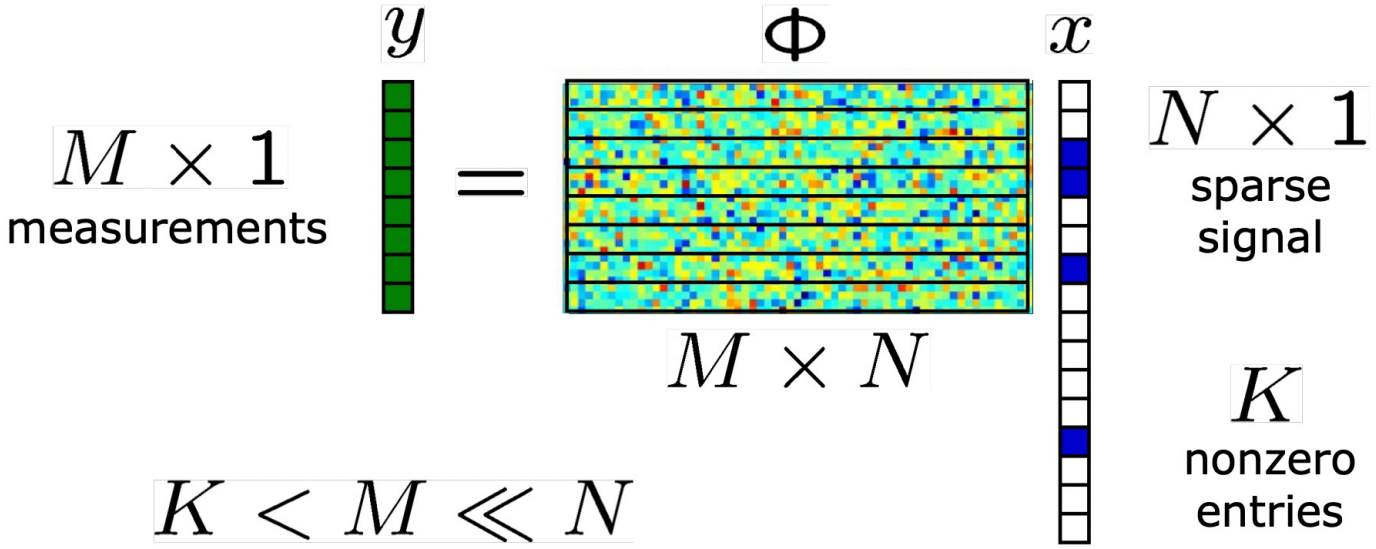
$$\alpha_i = \langle x, \psi_i \rangle$$



$K \ll N$
large
wavelet
coefficients

CS Signal Recovery

- Reconstruction/decoding: given $y = \Phi x$
  (ill-posed inverse problem) find x



- Reconstruction/decoding: given $y = \Phi x$
  (ill-posed inverse problem) find x

- L_2 fast, wrong

$$\hat{x} = \arg\min_{y=\Phi x} \|x\|_2$$

- L_0 **correct, slow**
  only $M=K+1$ measurements required to perfectly reconstruct K-sparse signal
  [Bresler; Wakin et al]

$$\hat{x} = \arg\min_{y=\Phi x} \|x\|_0$$

number of nonzero entries

- Reconstruction/decoding: given $y = \Phi x$
  (ill-posed inverse problem) find x

- L_2 fast, wrong $\quad \widehat{x} = \arg\min_{y=\Phi x} \|x\|_2$

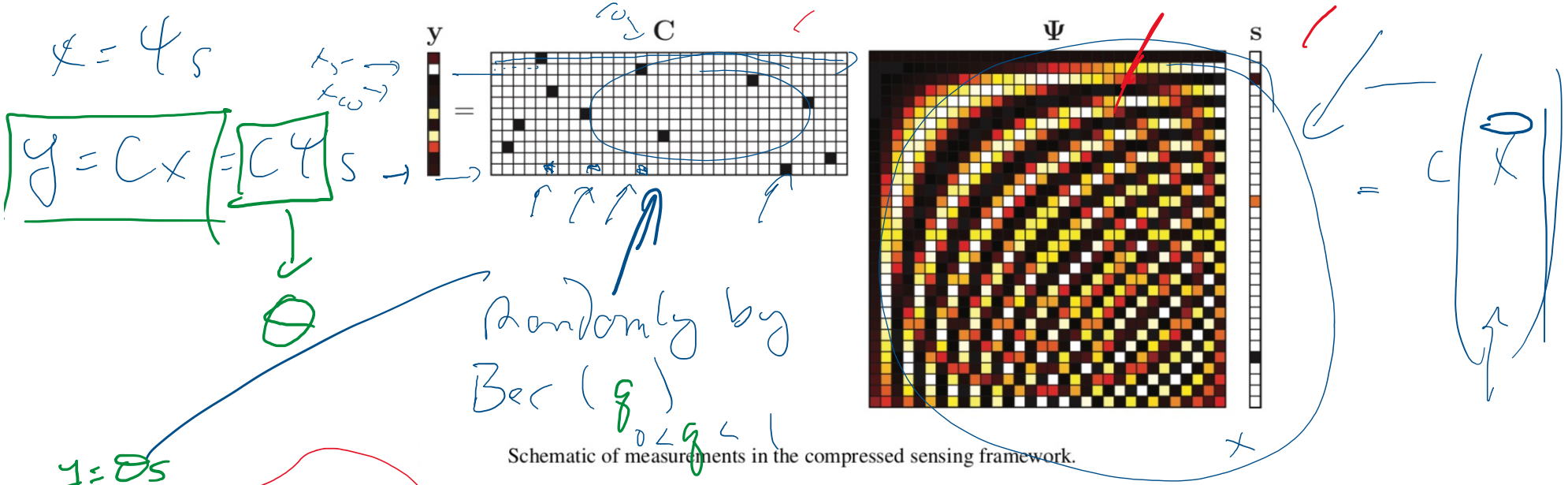
- L_0 correct, slow $\quad \widehat{x} = \arg\min_{y=\Phi x} \|x\|_0$

- L_1 **correct, mild oversampling** $\quad \widehat{x} = \arg\min_{y=\Phi x} \|x\|_1$

  [Candes et al, Donoho]

  linear program

$$M \approx K \log N \ll N$$



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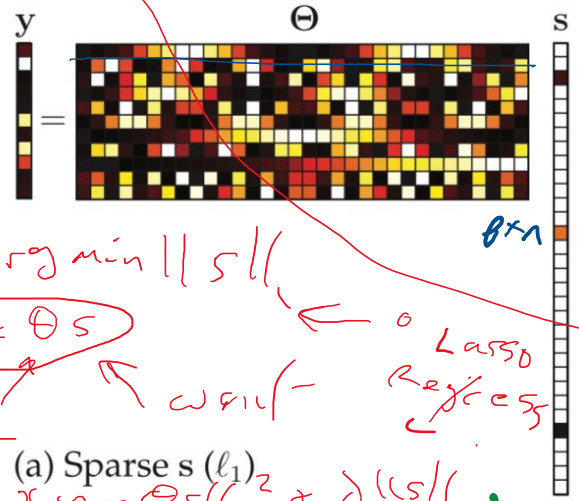
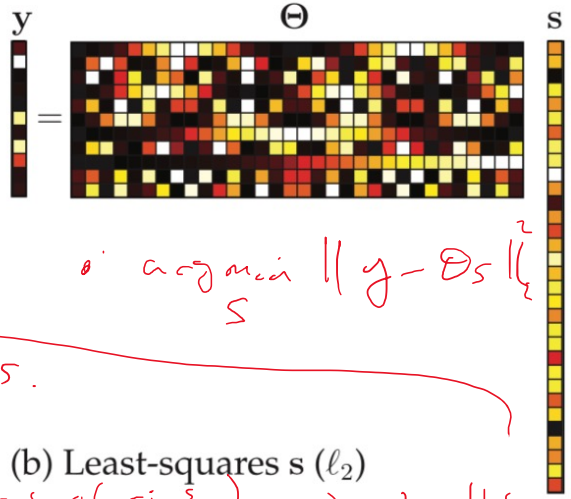


(a) Sparse s (ℓ_1)

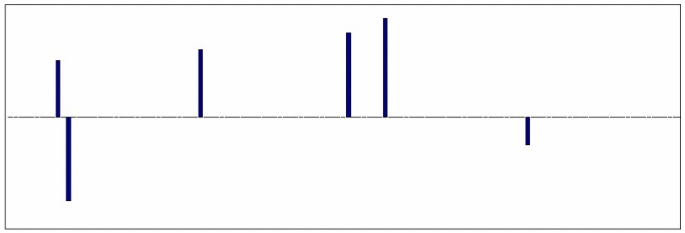


(b) Least-squares s (ℓ_2)

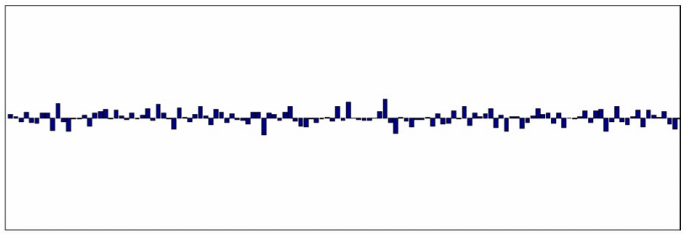
CS Signal Recovery

- Reconstruction/decoding: given $y = \Phi x$
  (ill-posed inverse problem) find x

- L_2 fast, **wrong** $\hat{x} = \arg\min_{y=\Phi x} \|x\|_2$



x

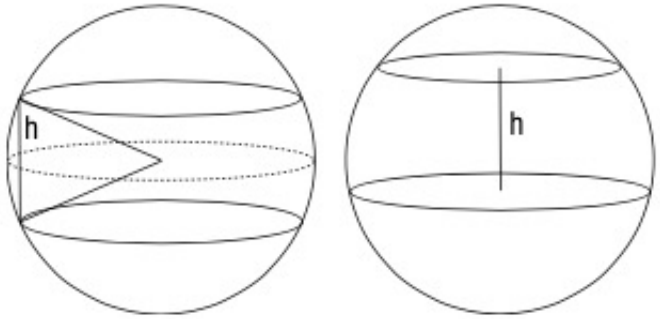


$\hat{x} = (\Phi^T \Phi)^{-1} \Phi^T y$

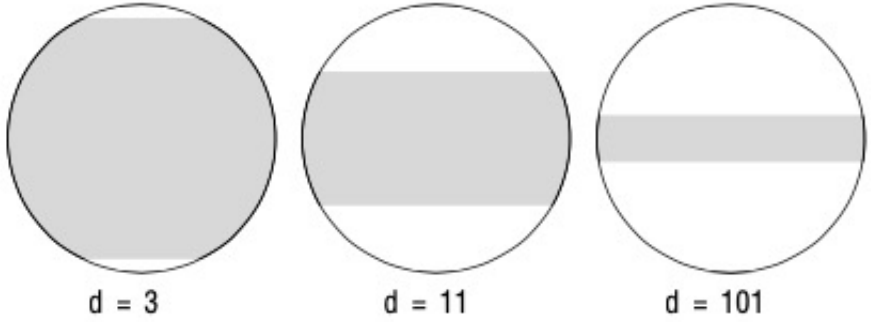
Another view on **where all the random goes** …geometry of data in high dimensions

highly important topic in Machine Learning, Data Science, and ROM

-story one – random gets soaked up by time.
-story two – its about random projection (Johnson-Lindenstrass) which is an application of concentration of measure



a strip of width h



(width of the) strip around the equator that contains 90% of the area

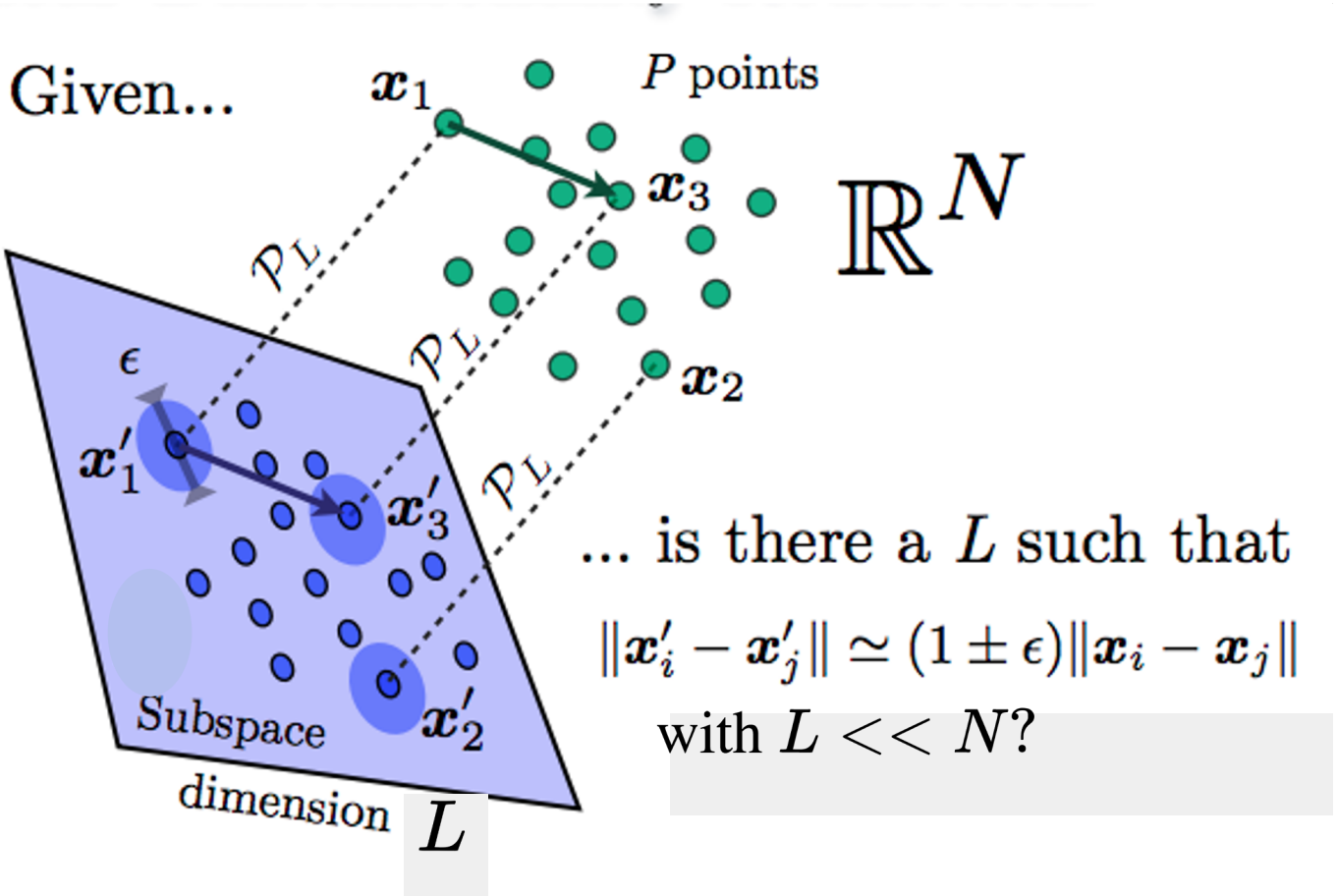
Several different important concepts follow from this simple idea
-Markov inequalities, Chernoff inequality

-for us – in high dimensions, samples of data ”look” almost like they live in a “flatter” space – Johnson-Lindenstrass theorem

See Ortega for CoM/JL random projection interpretation of RC

Matouš̌ek’s book [Mat02]

Random Projection



Source: https://perso.uclouvain.be/laurent.jacques/uploads/Main/buffon-present-120314-MLG.pdf

Random Projection – on projection and approximate isometry

- The random projection method is based on the Johnson-Lindenstrauss lemma.

Theorem[Johnson-Lindenstrauss Lemma]:

For any $0 < \epsilon < 1$ and any integer $M > 1$, let L be a positive integer such that $L \geq L_0$ with $L_0 = \frac{C \ln M}{\epsilon^2}$,

where C is a suitable constant ($C \approx 8$ in practice,$C = 2$ is good enough). Then for any set X of M data points in $\mathbb{R}^N$,

there exists a map $f : \mathbb{R}^N \to \mathbb{R}^L$ such that

for all $x_1, x_2 \in X, (1-\epsilon)||x_1 - x_2||^2 \leq ||f(x_1) - f(x_2)||^2 \leq (1+\epsilon)||x_1 - x_2||^2$.

[Johnson et. al, 1984]

Theorem[Random Projection]

For any $0 < \epsilon, \delta < \frac{1}{2}$ and positive integer N,

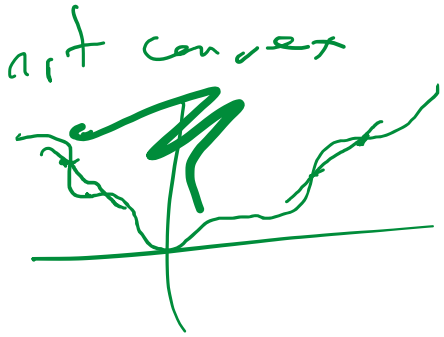
there exits a random matrix of B of size $L \times N$ such that for $L \geq L_0$ with $L_0 = \frac{C \ln(1/\delta)}{\epsilon^2}$

and for any unit-length vector $x \in R^N$, $Pr\{|||Bx||^2 - 1| > \epsilon\} \leq \delta$

or $Pr\{|||Bx||^2 - 1| > \epsilon\} \leq e^{-CL\epsilon^2}$

[Papadimitriou et. al., 1998]

CS and linprog



CVX examples

To solve

$$\min_x \|x\|_1 \quad Ax = b$$

$$\min_{x,t} \sum_{i=1}^{n} t_i \quad x \le t \quad -x \le t \quad Ax = b$$

Use CVX code

```
cvx_begin
  variable x(n);
  variable t(n);
  minimize(sum(t));
  subject to
    x <= t;
    -x <= t;
    A*x == b;
cvx_end
```